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12/30/2003

David Qiang Meng

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09/30/2011

MISSION/BSTZ

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**Please find below and/or attached an Office communication concerning this application or proceeding.**

The time period for reply, if any, is set in the attached communication.

UNITED STATES PATENT AND TRADEMARK OFFICE

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BEFORE THE BOARD OF PATENT APPEALS  
AND INTERFERENCES

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*Ex parte* DAVID QIANG MENG

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Appeal 2009-009769  
Application 10/750,423  
Technology Center 2100

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Before LANCE LEONARD BARRY, HOWARD B. BLANKENSHIP, and  
JAMES R. HUGHES, *Administrative Patent Judges*.

BLANKENSHIP, *Administrative Patent Judge*.

DECISION ON APPEAL

STATEMENT OF THE CASE

This is an appeal under 35 U.S.C. § 134(a) from the Examiner's final rejection of claims 1-26, which are all the claims in the application. We have jurisdiction under 35 U.S.C. § 6(b).

We affirm.

*Representative Claim*

1. A computer-implemented method comprising:

partitioning a memory device to produce a first group of memory entries being accessible in parallel and selectable independent of a second group of memory entries in the memory device that is accessible in parallel.

*Examiner's Rejections*

Claims 1-26 stand rejected under 35 U.S.C. § 112, first paragraph, as failing to comply with the enablement and the written description requirements.<sup>1</sup>

PRINCIPLES OF LAW

“Although the statute does not say so, enablement requires that the specification teach those in the art to make and use the invention without ‘undue experimentation.’ That some experimentation may be required is not fatal; the issue is whether the amount of experimentation required is ‘undue.’” *In re Vaeck*, 947 F.2d 488, 495 (Fed. Cir. 1991) (citations omitted).

“Whether undue experimentation is needed is not a single, simple factual determination, but rather is a conclusion reached by weighing many factual considerations.” *In re Wands*, 858 F.2d 731, 737 (Fed. Cir. 1988). The factors to be considered in determining whether a disclosure would require undue experimentation include:

- (1) the quantity of experimentation necessary, (2) the amount of direction or guidance presented, (3) the presence or absence of working examples, (4) the nature of the invention, (5) the state

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<sup>1</sup> The Examiner has withdrawn a § 101 utility rejection of the claims.

of the prior art, (6) the relative skill of those in the art, (7) the predictability or unpredictability of the art, and (8) the breadth of the claims.

*Id.*

The full scope of the claimed invention must be enabled. The rationale for this statutory requirement is straightforward. Enabling the full scope of each claim is “part of the quid pro quo of the patent bargain.” A patentee who chooses broad claim language must make sure the broad claims are fully enabled. “The scope of the claims must be less than or equal to the scope of the enablement” to “ensure[ ] that the public knowledge is enriched by the patent specification to a degree at least commensurate with the scope of the claims.”

*Strick v. Dreamworks*, 516 F.3d 993, 999 (Fed. Cir. 2008) (alteration in original) (citations omitted).

The examiner bears the initial burden of setting forth a reasonable explanation as to why the scope of protection provided by the claims is thought to be not adequately enabled by the description of the invention provided in the specification. If that burden is met, the burden then shifts to the applicant to provide proof that the specification is indeed enabling. *In re Wright*, 999 F.2d 1557, 1561-62 (Fed. Cir. 1993).

## ANALYSIS

### *Section 112 rejections*

In response to the § 112, first paragraph rejection for lack of enablement, Appellant argues claim 1 (App. Br. 8-11)<sup>2</sup>. Appellant also

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<sup>2</sup> In making our determinations, we have considered Appellant’s disclosure, the Final Rejection (mailed Oct. 18, 2006), the Appeal Brief (filed June 18,

presents nominal arguments with respect to dependent claims 2 through 7 and places each of the dependent claims in different headings (*id.* 12-13). However, the remarks for the dependent claims at best merely refer to alleged support in the Specification and allege that the claims are enabled without explaining how the alleged support might serve to enable the claims. Further, the remarks do not allege that the limitations of the dependent claims would remedy any deficiency in the enablement of base claim 1. The remarks with respect to the written description rejection (*id.* 13-14) allege in general that the claims are supported by the disclosure (including the original claims), but do not point to support in the disclosure for any particular claim. Accordingly, we will decide the appeal on the basis of claim 1 alone. *See* 37 C.F.R. § 41.37(c)(1)(vii).

The Examiner refers to Handy as evidence of what constitutes a conventional CAM (content addressable memory). As depicted and described in Figure 1.7 of the reference, in a CAM an address presented to the compare address bus is compared simultaneously with the contents of every memory location. If there is a match at any location, the address of that location is sent to the N-bit binary output pins. The Examiner finds that the CAM shown in Handy does not appear to have the ability to be partitioned into two groups of entries as recited in the claims. Ans. 16.

Appellant's Specification provides:

Each of the entries in CAM 54 [Fig. 3] is configurable by a CAM manager 58 that is implemented as microcode in the control store 50 and, which is executed by the packet engine 48.

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2007), the Examiner's Answer (mailed Sept. 18, 2007), the Reply Brief (filed Nov. 19, 2007), the Supplemental Examiner's Answer (mailed Feb. 1, 2008), and the Supplemental Reply Brief (filed Mar. 17, 2008).

The CAM manager 58 partitions the CAM 54 into a particular number of entries [Fig 3]. The CAM manager 58 is capable of partitioning individual entries into two or more subentries that are individually selectable for use in parallel comparisons. By producing subentries, particular ones of the subentries are grouped for storing one type of data (e.g. MAC addresses) and selected for use in comparing the data in parallel. Other subentries in the same CAM entries are grouped for storing and comparing another type of data (e.g., IP addresses). [Fig. 4]. Thus, CAM 54 is configured by CAM manager 58 for storing two or more types of data in subentries that are individually selectable for use in parallel comparisons.

Spec. 9:18 - 10:8 (references to Figures inserted).

The description for producing what is purported to be a novel, non-obvious type of CAM seems to be contained in the words “is configurable by a CAM manager,” which is “implemented as microcode” in a control store and is executed by a “packet engine.” Appellant’s Specification offers no explanation of how the CAM manager actually configures, i.e., makes, the number of CAM entries and subentries.

While the written description and drawings provide further examples of differently configured CAMs, we agree with the Examiner to the extent that the disclosure seems to provide little beyond describing new rectangles and new lettering in the drawings that represent new entries or subentries. The disclosure does not teach how to make and use the newly configured CAMs.

Further, while we agree with Appellant to the extent that the electrical arts are considered to be relatively predictable and that the skill in the pertinent art appears to be relatively high, those facts are outweighed by the facts that tend to show that undue experimentation would be required. The

dearth of direction or guidance presented and the absence of working examples weigh against Appellant's position. Further, instant claim 1 is not limited to partitioning a CAM (which Appellant contends the disclosure enables), but by its terms extends to any “memory device.” The written description and drawings do not purport to teach how to partition *any* memory device to produce groups of memory entries that are accessible in parallel and independently selectable.

We conclude that the Examiner has set forth a reasonable explanation as to why the scope of protection provided by the claims is thought to be not adequately enabled. Appellant has not produced evidence in rebuttal to call the conclusion of non-enablement into question.<sup>3</sup>

We sustain the § 112 rejection for lack of enablement. We also sustain the § 112 rejection for lack of written description. The Examiner and Appellant seem to acknowledge that the rejection for lack of written description stands or falls with the rejection for lack of enablement. “Those two requirements usually rise and fall together. That is, a recitation of how to make and use the invention across the full breadth of the claim is ordinarily sufficient to demonstrate that the inventor possesses the full scope of the invention, and vice versa.” *LizardTech, Inc. v. Earth Res. Mapping*, 424 F.3d 1336, 1345 (Fed. Cir. 2005).

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<sup>3</sup> We note that Appellant refers in the reply briefs to evidence that is not before us, which was apparently filed with an after-appeal Information Disclosure Statement. *See* 37 C.F.R. § 41.37(c)(1)(ix) (*Evidence appendix*). In any event, to the extent the evidence is presented in the reply briefs, the evidence would not constitute proof that the disclosure is enabling.

*Prior Art*

A cursory search of the US patent database yielded US 6,324,087 B1 (Pereira), which appears to anticipate instant claim 1. Pereira describes method and apparatus for partitioning a content addressable memory device (Title). As summarized at column 1, line 43 through column 2, line 6, and further detailed at column 3, lines 5 through 12 and column 4, lines 5 through 29, the reference describes a method for partitioning a CAM device (Fig. 1) having a plurality of CAM blocks into a number of individually searchable partitions. Block select circuits 106(1) through 106(n) selectively disable corresponding CAM blocks from participating in the compare operation. As Appellant admits (Reply Br. 11), since a CAM searches all entries (simultaneously), it automatically searches in parallel, at least to the level of the subject matter of the claims.

That prior art may anticipate instant claim 1 does not mean that the instant disclosure is enabling for the subject matter thereof. First, the instant disclosure does not hint that the partitioning is accomplished by means of block select circuits as described by Pereira; apparently, the partitioning is achieved in some other way that is not adequately described. Second, the disclosure, at the least, fails to enable the full scope of the claimed invention. *See Strick v. Dreamworks*, 516 F.3d at 999. That instant claim 1 may be so broad that a reference falls within its scope does not necessarily mean that the disclosure meets the enablement requirement, but underscores a claim breadth that weighs toward the necessity for undue experimentation.



Appeal 2009-009769  
Application 10/750,423

DECISION

The rejections of claims 1-26 under 35 U.S.C. § 112, first paragraph, as failing to comply with the enablement and the written description requirements are affirmed.

AFFIRMED

pgc

*Notice of References Cited*

Application/Control No.

10/750,423

Applicant(s)/Patent Under  
Reexamination

Examiner

Jared Rutz

Art Unit

2800

Page 1 of 1

## U.S. PATENT DOCUMENTS

*		Document Number Country Code-Number-Kind Code	Date MM-YYYY	Name	Classification	
X	A	US-6,324,087	11-2001	Pereira	--	--
	B	US-				
	C	US-				
	D	US-				
	E	US-				
	F	US-				
	G	US-				
	H	US-				
	I	US-				
	J	US-				
	K	US-				
	L	US-				
	M	US-				

## FOREIGN PATENT DOCUMENTS

*		Document Number Country Code-Number-Kind Code	Date MM-YYYY	Country	Name	Classification	
	N						
	O						
	P						
	Q						
	R						
	S						
	T						

## NON-PATENT DOCUMENTS

*		Include as applicable: Author, Title Date, Publisher, Edition or Volume, Pertinent Pages)
	U	
	V	
	W	
	X	

\*A copy of this reference is not being furnished with this Office action. (See MPEP § 707.05(a).)  
 Dates in MM-YYYY format are publication dates. Classifications may be US or foreign.



US006324087B1

(12) **United States Patent**  
**Pereira**

(10) **Patent No.:** **US 6,324,087 B1**  
(45) **Date of Patent:** **Nov. 27, 2001**

(54) **METHOD AND APPARATUS FOR  
PARTITIONING A CONTENT ADDRESSABLE  
MEMORY DEVICE**

(75) Inventor: **Jose Pio Pereira**, Santa Clara, CA (US)

(73) Assignee: **NetLogic Microsystems, Inc.**,  
Mountain View, CA (US)

(\*) Notice: Subject to any disclaimer, the term of this  
patent is extended or adjusted under 35  
U.S.C. 154(b) by 0 days.

(21) Appl. No.: **09/590,642**

(22) Filed: **Jun. 8, 2000**

(51) **Int. Cl.**<sup>7</sup> ..... **G11C 15/00**

(52) **U.S. Cl.** ..... **365/49; 365/189.07; 365/230.03**

(58) **Field of Search** ..... 365/49, 78, 190,  
365/189.07, 230.03

(56) **References Cited**

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0 872 802	10/1998	(EP)

\* cited by examiner

*Primary Examiner*—David Nelms

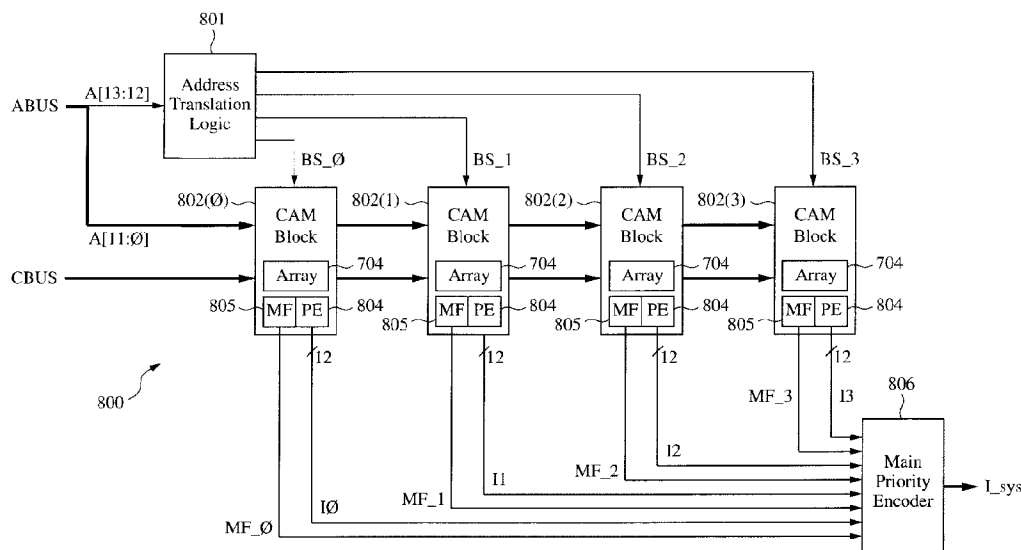
*Assistant Examiner*—Thong Le

(74) *Attorney, Agent, or Firm*—William L. Paradise

(57) **ABSTRACT**

A CAM device having a plurality of CAM blocks is partitioned into a number of individually searchable partitions, where each partition may include one or more CAM blocks of the CAM device. In one embodiment, each CAM block is connected to a block select circuit that stores a class code indicating what class or type of data is stored in the block. The same class code may be stored in any number of block select circuits to define a partition as including the corresponding number of CAM blocks. During compare operations between a comparand word and data stored in the CAM device, a search code is provided to the block select circuits. Each block select circuit compares the search code with its class code and, in response thereto, selectively enables or disables the corresponding CAM block for the compare operation. In some embodiments, the block select circuit enables the corresponding CAM block if the search class matches the class code and, conversely, disables the corresponding CAM block if the search code does not match the class code.

**33 Claims, 11 Drawing Sheets**



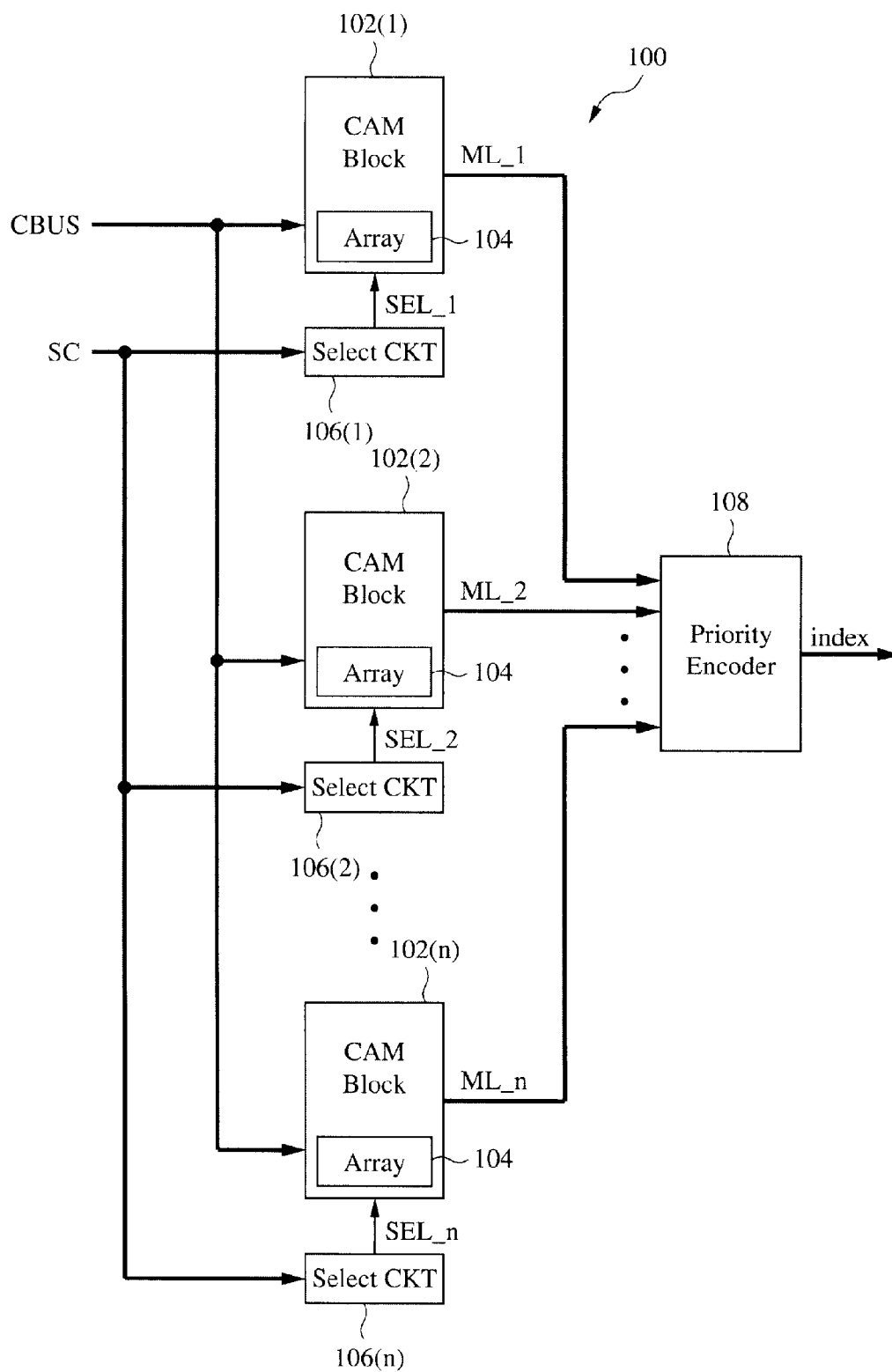


FIG. 1

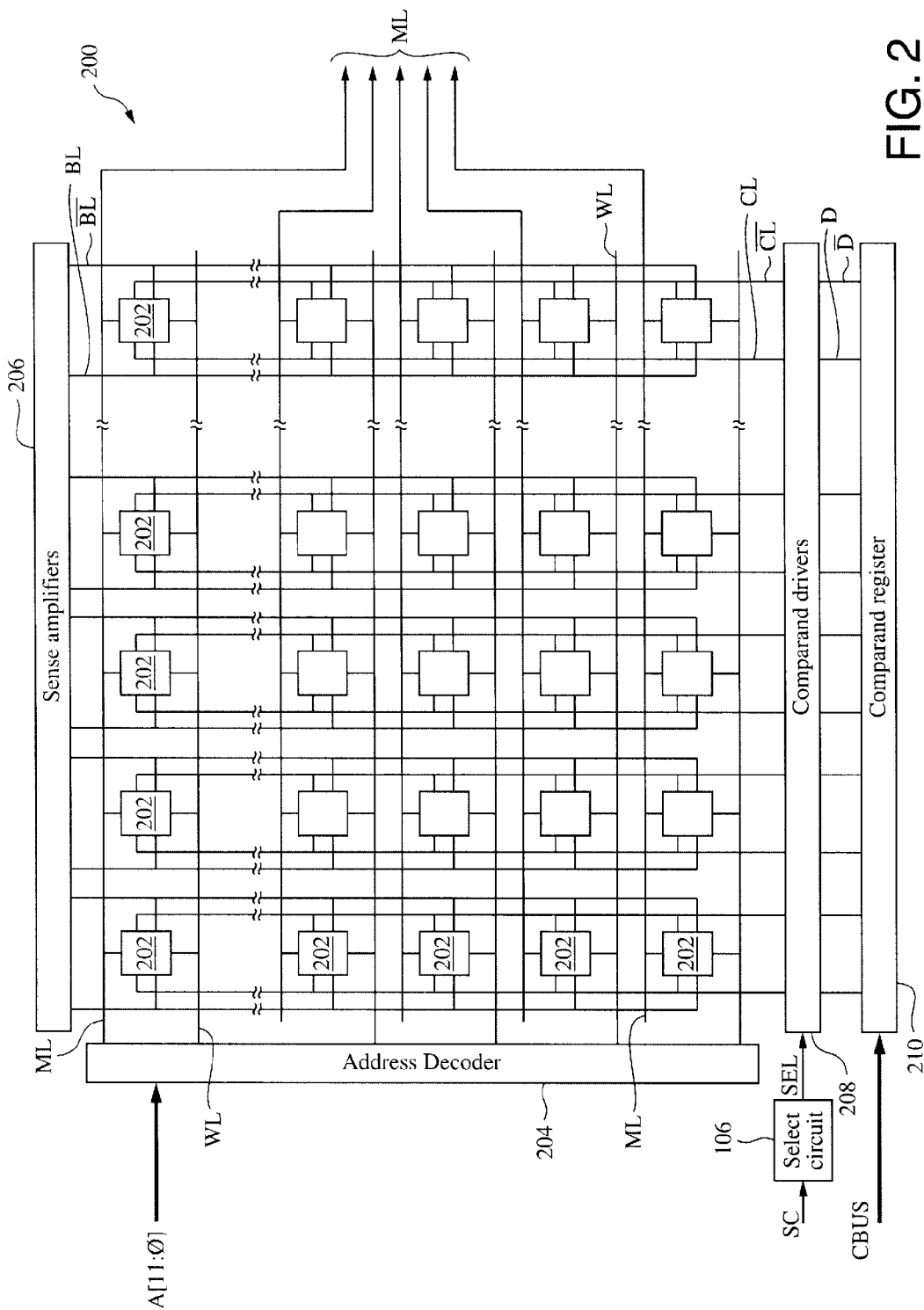


FIG. 2

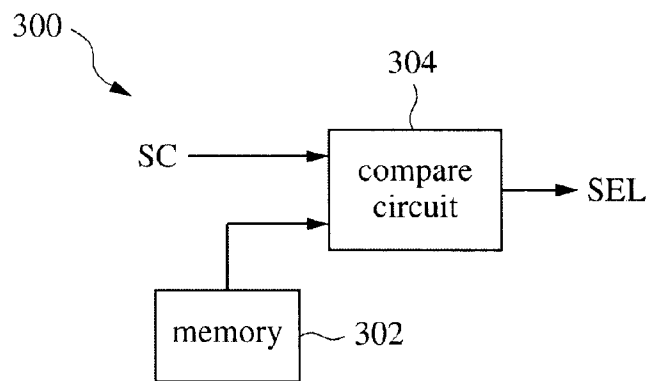


FIG. 3

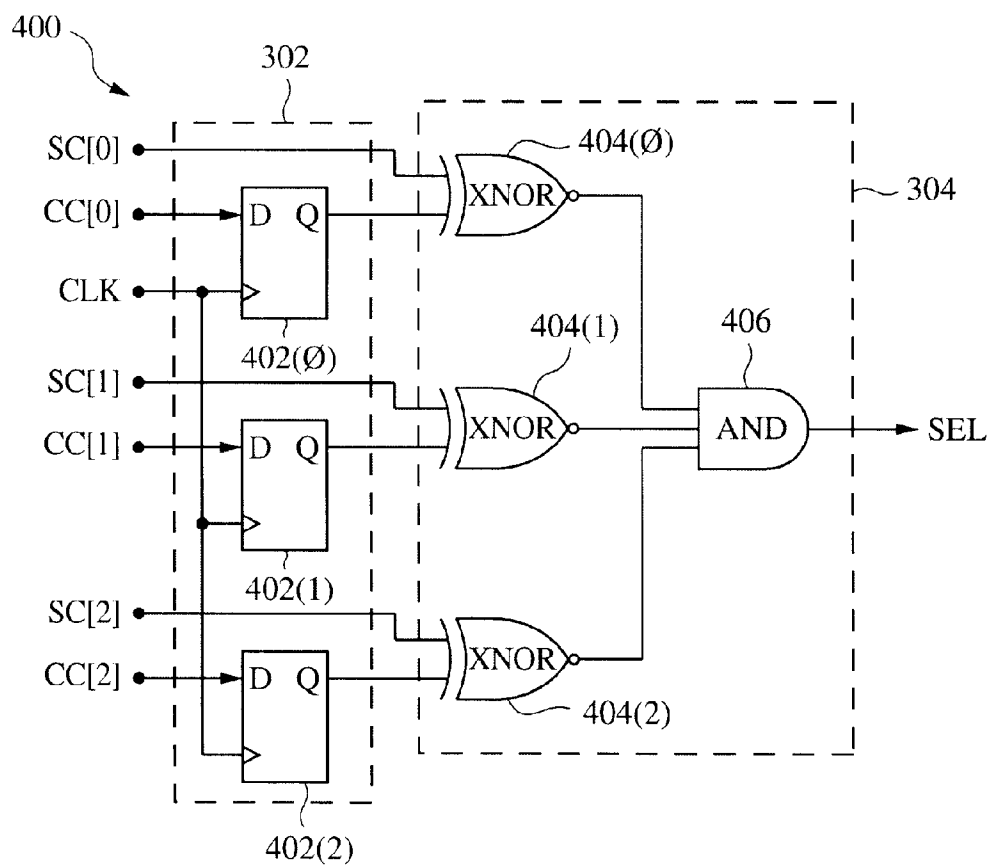


FIG. 4

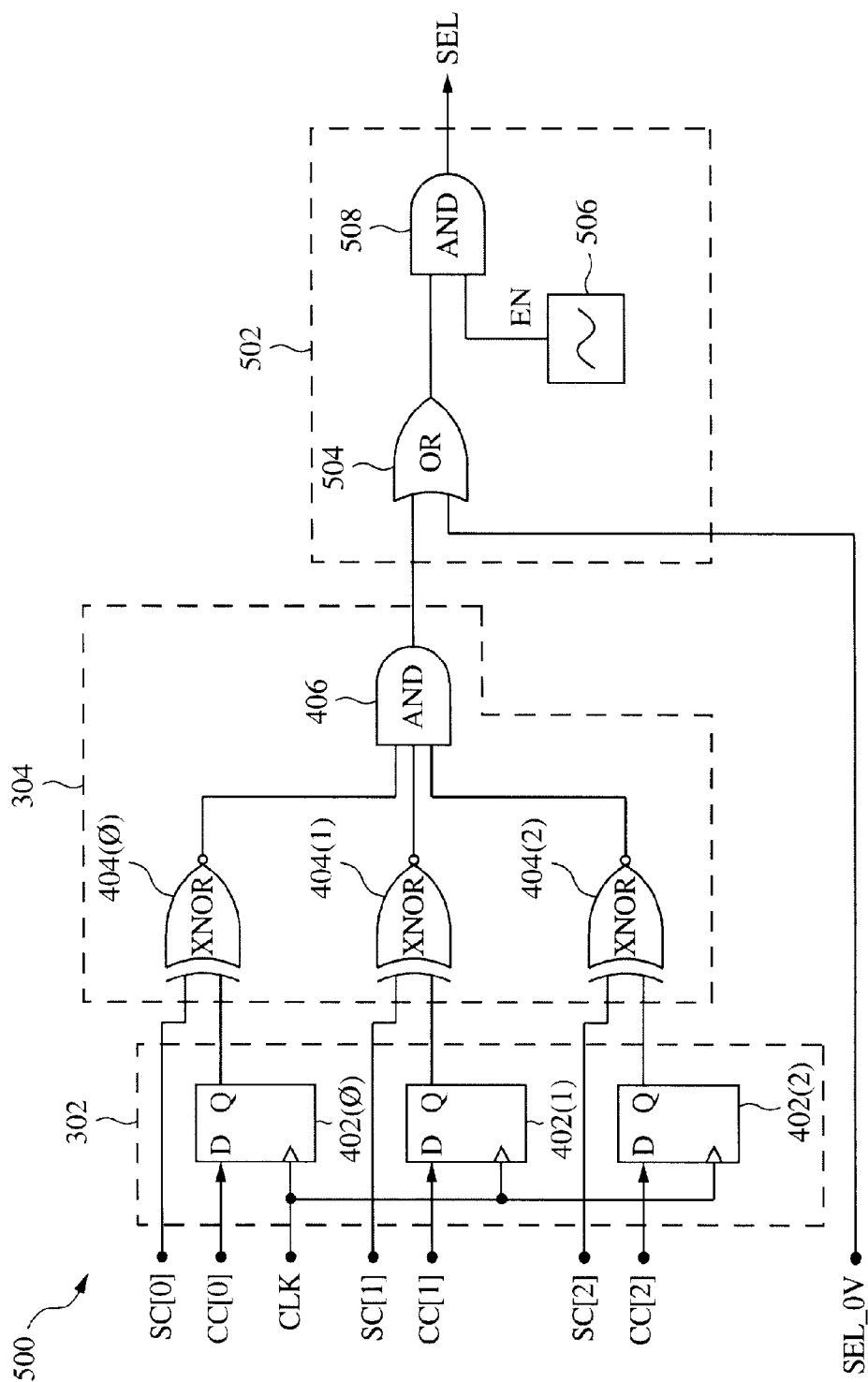


FIG. 5

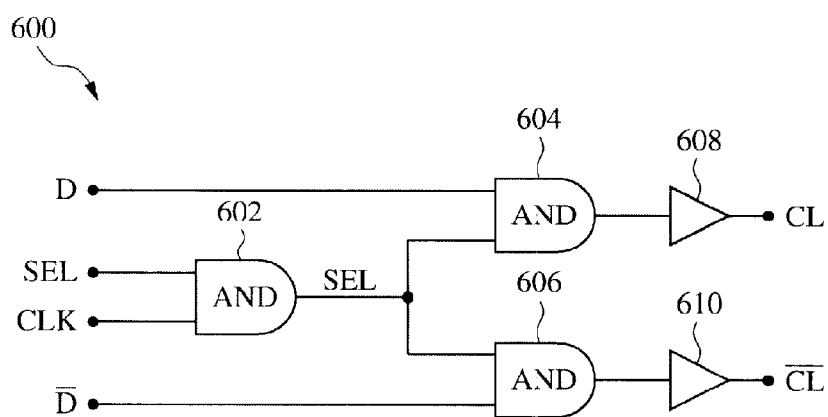


FIG. 6



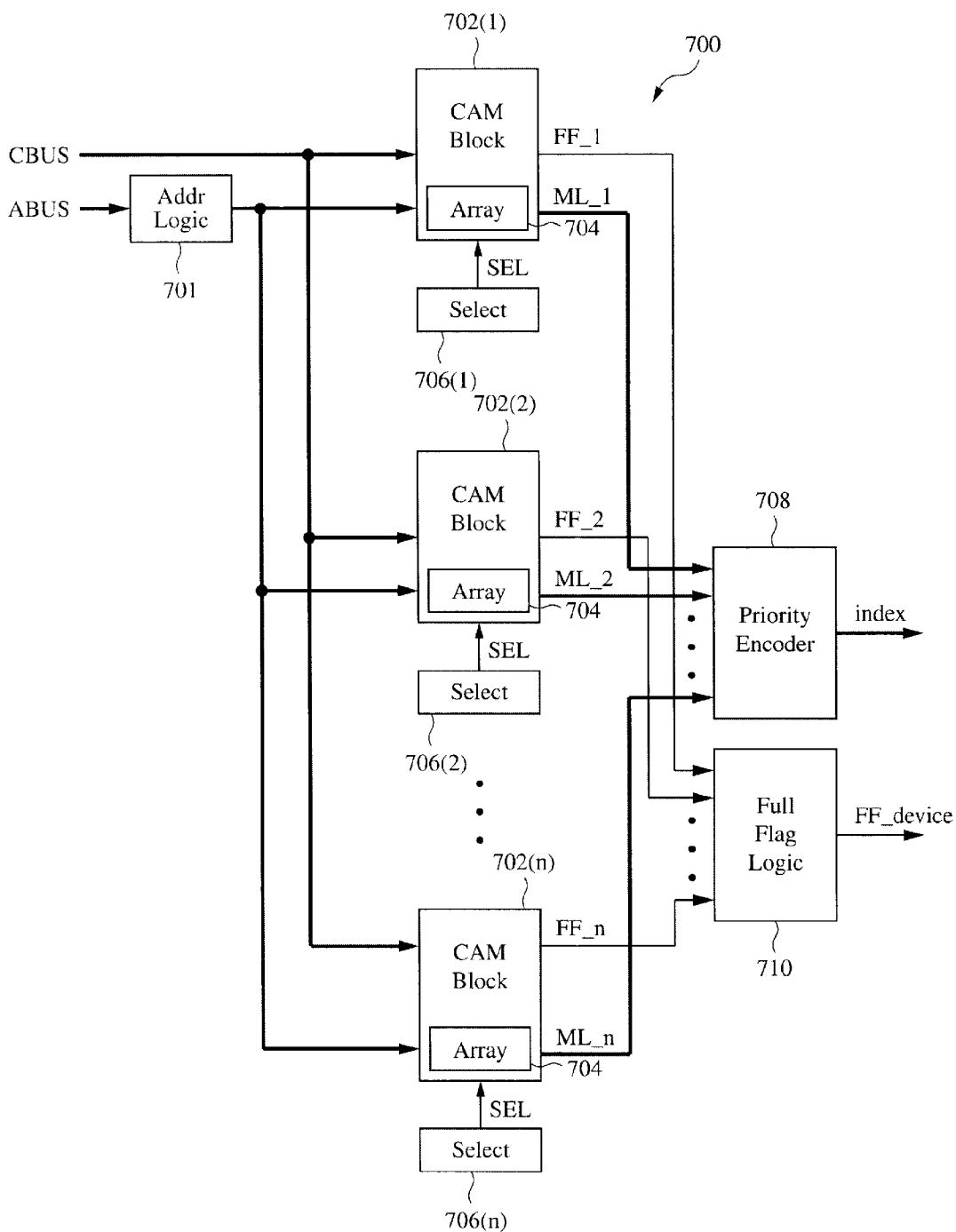


FIG. 7

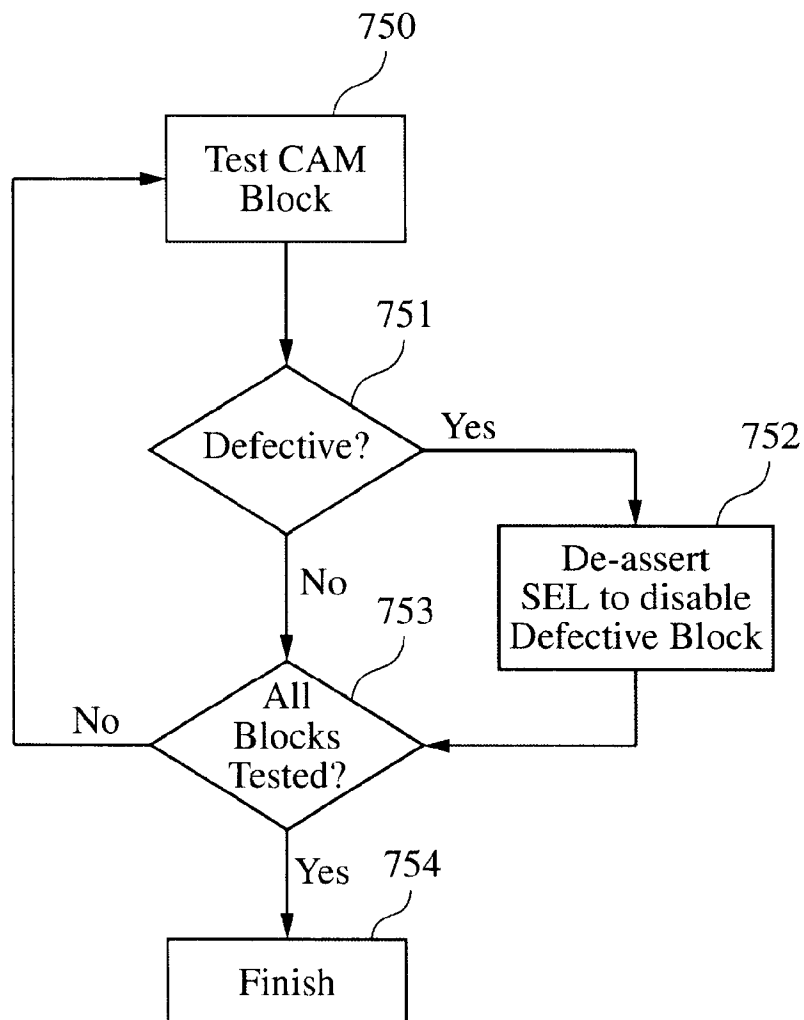


FIG. 8

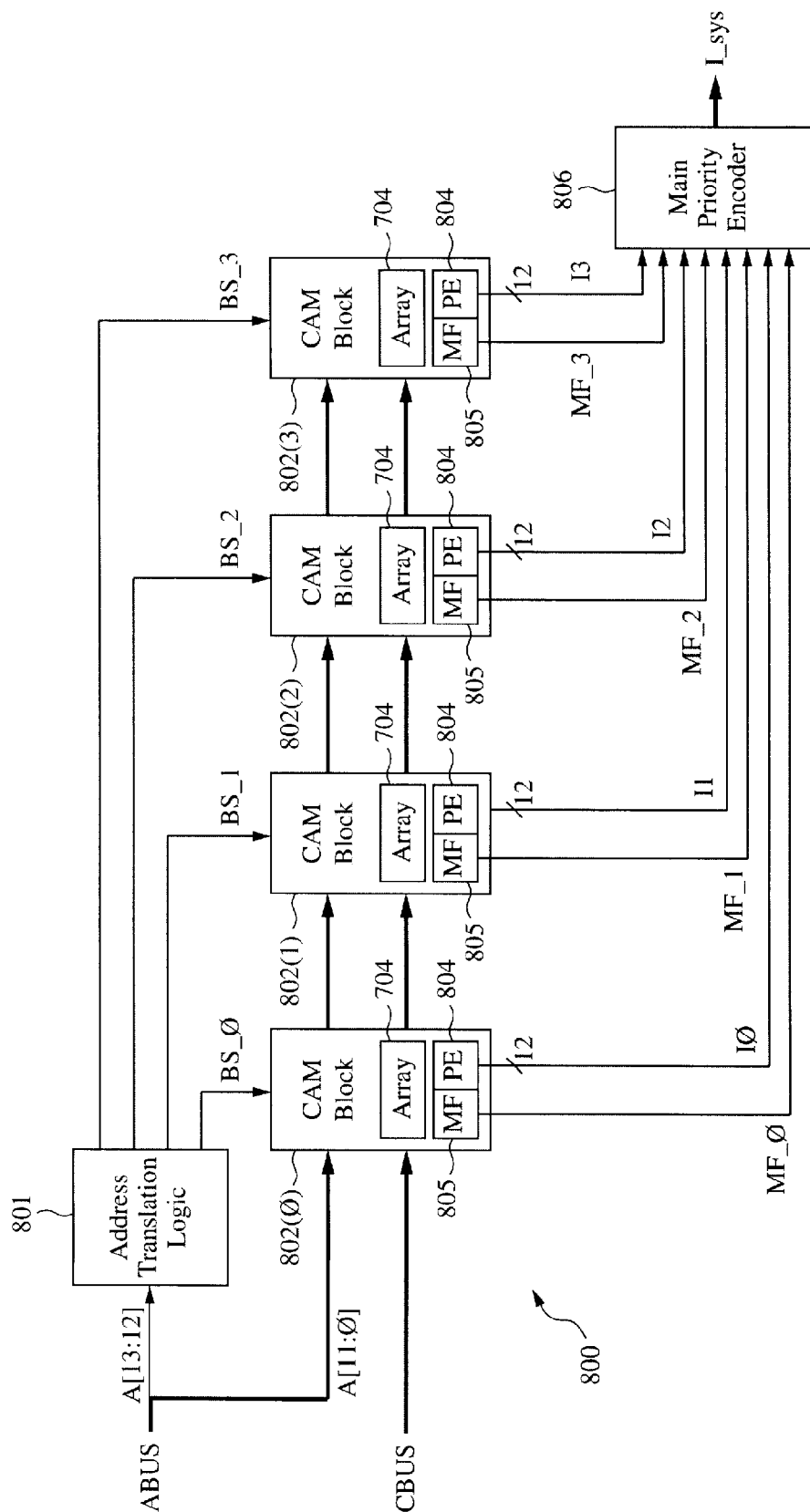


FIG. 9

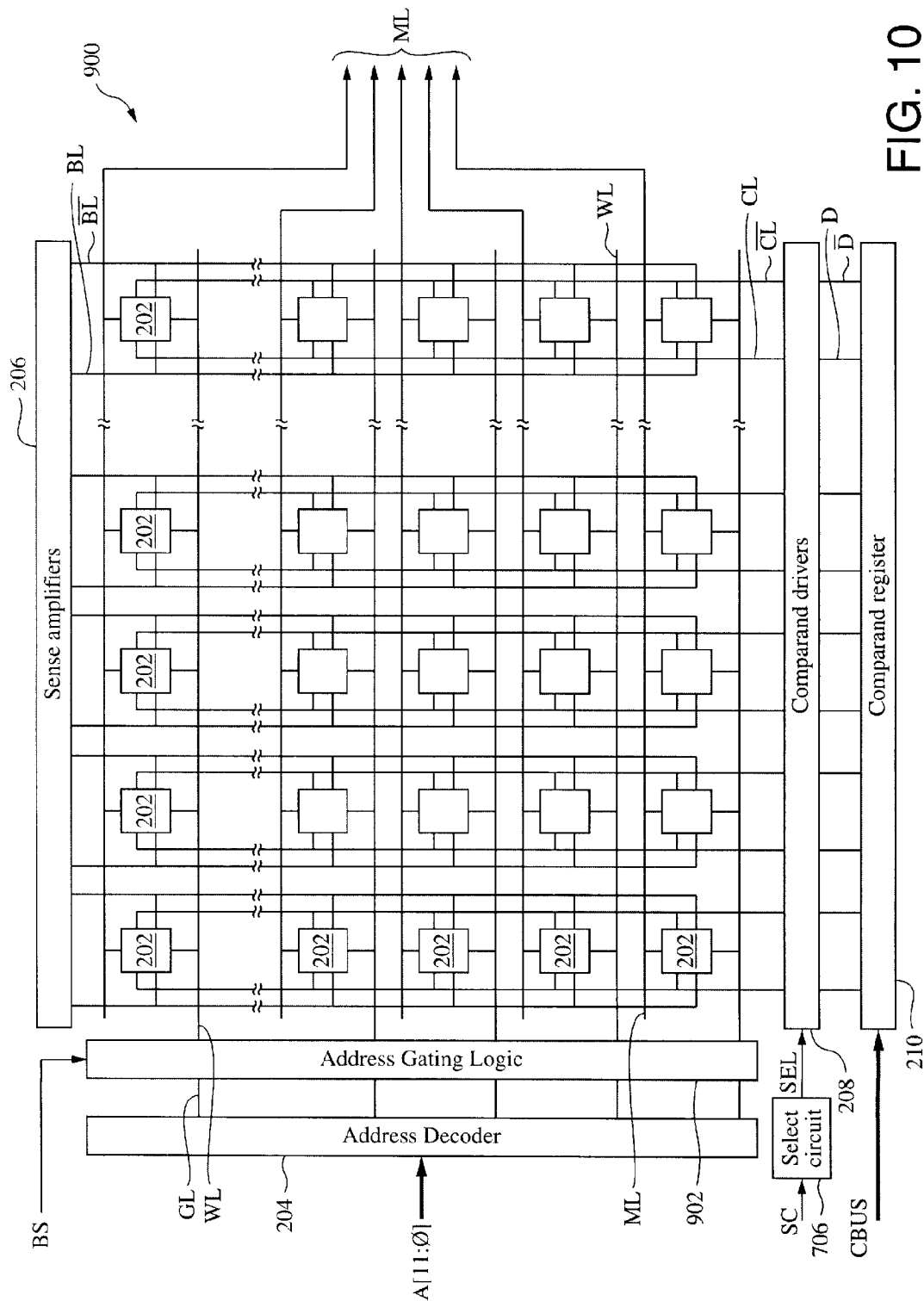


FIG. 10

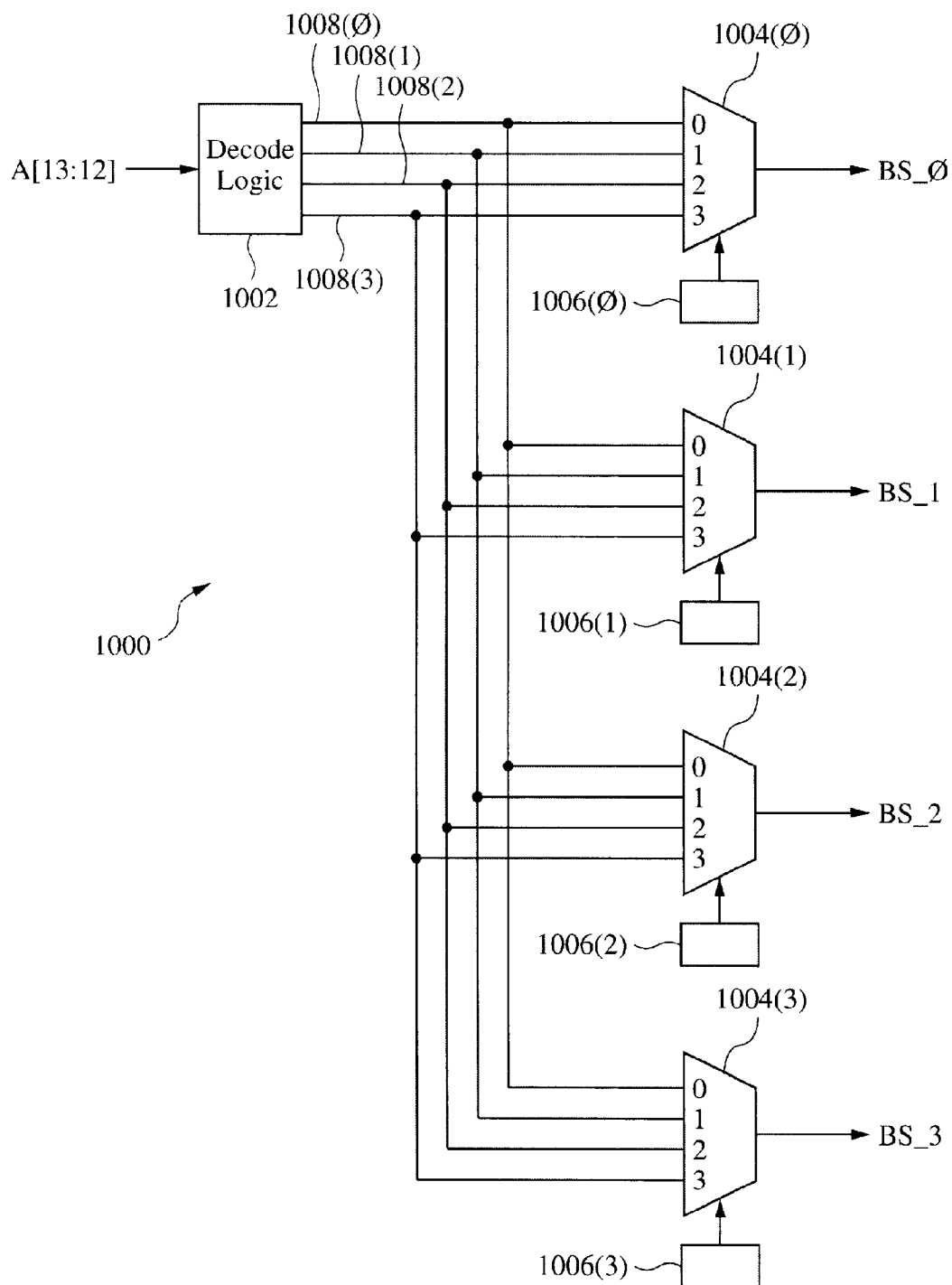


FIG. 11

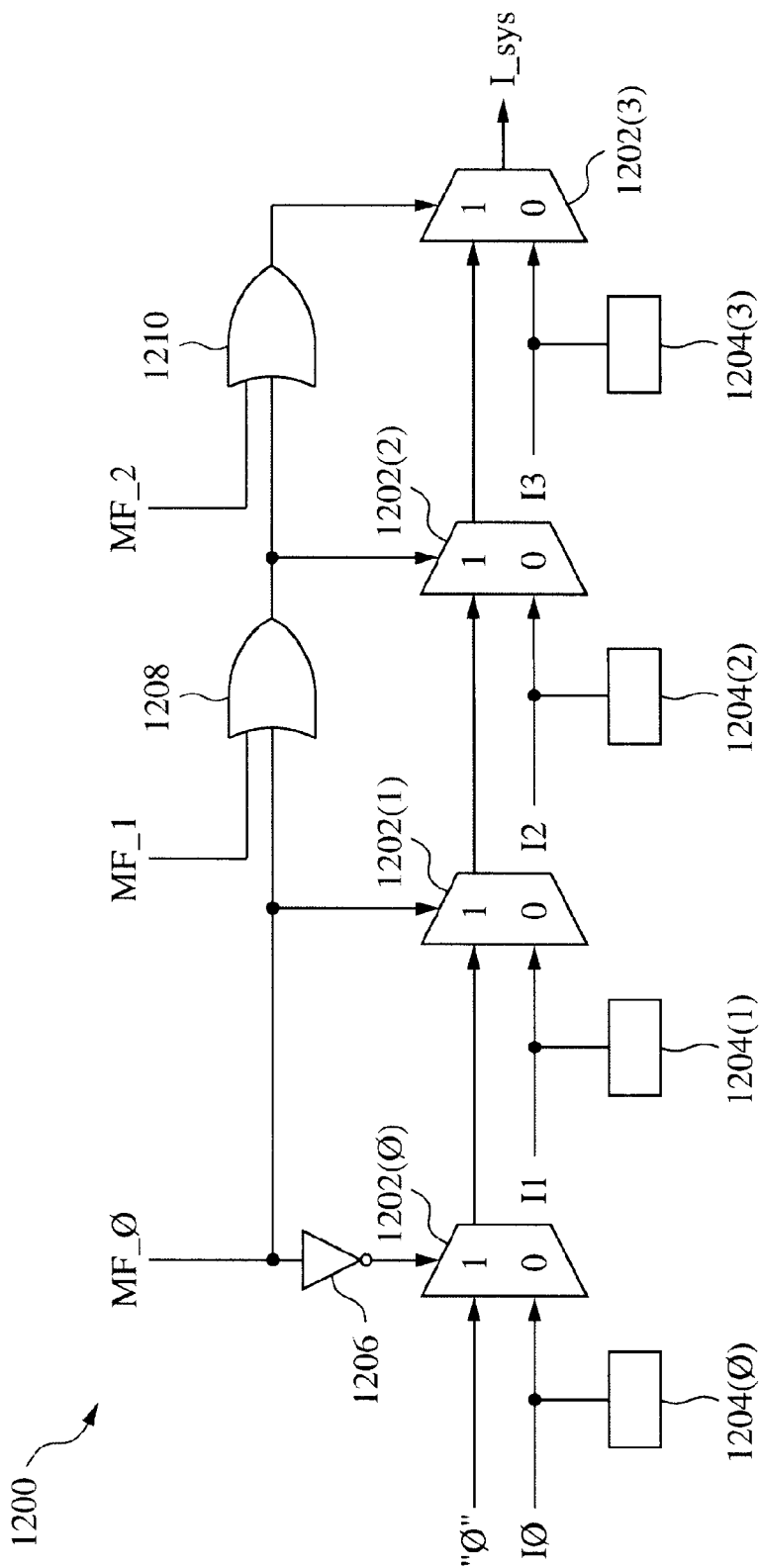


FIG. 12

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# METHOD AND APPARATUS FOR PARTITIONING A CONTENT ADDRESSABLE MEMORY DEVICE

## BACKGROUND

### 1. Field of Invention

This invention relates generally to semiconductor memories and specifically to content addressable memories.

### 2. Description of Related Art

Content addressable memories (CAMs) are frequently used for address look-up functions in Internet data routing. For example, routers used by local Internet Service Providers (ISPs) typically include one or more CAMs for storing a plurality of Internet addresses and associated data such as, for instance, corresponding address routing information. When data is routed to a destination address, the destination address is compared with all CAM words, e.g., Internet addresses, stored in the CAM array. If there is a match, routing information corresponding to the matching CAM word is output and thereafter used to route the data.

A CAM device includes a CAM array having a plurality of memory cells arranged in an array of rows and columns. Each memory cell stores a single bit of digital information, i.e., either logic zero or logic one. The bits stored within a row of memory cells constitute a CAM word. During compare operations, a comparand word is received at appropriate input terminals of a CAM device and driven into the CAM array using comparand lines to be compared with all the CAM words in the device. For each CAM word that matches the comparand word, a corresponding match line signal is asserted to indicate a match condition. If the comparand word matches more than one of the CAM words, the match line corresponding to each of the matching CAM words is asserted, and a "multiple match" flag is also asserted to indicate the multiple match condition. The match line signals from each CAM block are combined in a priority encoder to determine the index or address of the highest-priority matching CAM word. Associative information corresponding to the highest-priority matching CAM word stored in, for instance, an associated RAM, may also be provided.

A single CAM device can be used to store multiple tables each storing and maintaining different classes of data. All entries, however, typically participate in a compare operation. This can cause an undesirable amount of power to be drawn during the compare operation. It would be desirable to limit a search to only those entries associated with a particular class of data to reduce power consumption during the operation.

## SUMMARY

A method and apparatus are disclosed that may be used to partition a CAM device having a plurality of CAM blocks into a number of individually searchable partitions, where each partition may include one or more CAM blocks of the CAM device. In accordance with one embodiment of the present invention, each CAM block is connected to a block select circuit that stores a class code indicating what class or type of data is stored in the block. The same class code may be stored in any number of the block select circuits to define a partition as including the corresponding number of CAM blocks. During compare operations between a comparand word and data stored in the CAM device, a search code is provided to the block select circuits. Each block select circuit compares the search code with its class code and, in

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response thereto, selectively enables or disables the corresponding CAM block for the compare operation. In some embodiments, the block select circuit enables the corresponding CAM block if the search class matches the class code and, conversely, disables the corresponding CAM block if the search code does not match the class code.

In one embodiment, the block select circuit disables a corresponding CAM block by driving the comparand lines of the CAM block to a predetermined state to preclude the comparand word from being driven onto the comparand lines during the compare operation. By driving the comparand word only on the comparand lines of the selected (i.e., enabled) CAM blocks during the compare operation, present embodiments not only allow for selective searching across CAM blocks according to class codes, but also reduce power consumption in un-selected (i.e., disabled) CAM blocks during such selective compare operations.

## BRIEF DESCRIPTION OF THE DRAWINGS

FIG. 1 is a block diagram of a CAM device including a plurality of CAM blocks in accordance with one embodiment of the present invention;

FIG. 2 is a block diagram of a CAM block in one embodiment of the array of FIG. 1;

FIG. 3 is a block diagram of a block select circuit in one embodiment of the CAM block of FIG. 2;

FIG. 4 is a logic diagram of one embodiment of the block select circuit of FIG. 3;

FIG. 5 is a logic diagram of another embodiment of the block select circuit of FIG. 3;

FIG. 6 is a logic diagram of one embodiment of a comparand driver of the CAM block of FIG. 2;

FIG. 7 is a block diagram of a CAM device including a plurality of CAM blocks in accordance with another embodiment of the present invention configured to translate addresses of one or more defective CAM blocks;

FIG. 8 is a flow chart illustrating the disabling of defective CAM blocks in one embodiment;

FIG. 9 is a block diagram of one embodiment of the CAM device of FIG. 7;

FIG. 10 is a block diagram of a CAM block in one embodiment of the device of FIG. 9;

FIG. 11 is a block diagram of address translation logic in one embodiment of the CAM block of FIG. 10;

FIG. 12 is a logic diagram of a main priority encoder in one embodiment of the device of FIG. 9.

Like reference numerals refer to corresponding parts throughout the drawing figures.

## DETAILED DESCRIPTION

Embodiments of the present invention are discussed below in the context of a CAM device **100** for simplicity only. It is to be understood that embodiments of the present invention are equally applicable to CAM structures having other configurations of any suitable type of CAM cells. Further, architectural configurations of the present invention may be implemented in other types of memory blocks such as, for instance, RAM, Flash, and EEPROM. The interconnection between circuit elements or blocks may be shown as buses or as single signal lines, where each of the buses may alternatively be a single signal line, and each of the single signal lines may alternatively be a bus. In addition, the logic levels assigned to various signals in the description below are arbitrary, and therefore may be modified (e.g., reversed

polarity) as desired. Accordingly, the present invention is not to be construed as limited to specific examples described herein but rather includes within its scope all embodiments defined by the appended claims.

FIG. 1 shows a CAM device **100** in accordance with one embodiment of the present invention as having a number *n* of CAM blocks **102(1)–102(n)**, a corresponding number *n* of block select circuits **106(1)–106(n)**, and a priority encoder **108**. Each CAM block **102** includes a CAM array **104** having a plurality of rows of CAM cells for storing a plurality of CAM words therein, and is connected to a corresponding block select circuit **106**. Each row may also include one or more valid bits indicative of whether a valid CAM word is stored in the row. The valid bits may be used in a well-known manner to generate a full flag for the CAM block **102**. CAM blocks **102** may be any suitable type of CAM block, including for example, synchronous, asynchronous, binary, and ternary CAMs. Further, each CAM block **102** may be any suitable size, and in some embodiments may be of different sizes. In one embodiment, each CAM block **102** includes 1k (1024) rows of CAM cells.

During a compare operation, each CAM block **102** receives comparand data from a comparand bus CBUS. Other signals provided to the CAM device **100** during the compare operation may be a clock signal CLK, one or more instructions from an instruction decoder (not shown for simplicity), and other control signals. In some embodiments, instructions and comparand data may be provided to the CAM blocks **102(1)–102(n)** via the same bus. Other well-known signals which may be provided to the CAM blocks **102**, such as word enable signals, reset signals, and enable signals, are not shown for simplicity.

Each CAM block **102** provides a plurality of match line signals to the priority encoder **108** via corresponding match lines ML. The match lines carry match signals indicative of match conditions in the CAM arrays **104**. For simplicity, the plurality of match lines ML from each CAM block **102** are represented collectively in FIG. 1. The priority encoder **108** generates an index corresponding to one of the matching CAM words in the device **100**. In one embodiment, the priority encoder **108** outputs the index of the highest priority match. The highest priority match may be the lowest numbered address, the highest numbered address, or any other selected address.

For purposes of discussion herein, the first CAM block **102(1)** in the device **100** is designated as the highest priority block, the second CAM block **102(2)** is designated as the next highest priority block, and so on, and the last CAM block **102(n)** is designated as the lowest priority block, although in actual embodiments priority may be reversed or otherwise modified. Thus, the highest priority CAM block **102(1)** may include the lowest CAM addresses (i.e., CAM addresses 0 to *k*–1), the next highest priority block **102(2)** may include the next lowest CAM addresses (i.e., CAM addresses *k* to 2*k*–1), and so on, and the lowest priority CAM block **102(n)** may include the highest CAM addresses (i.e., CAM addresses (*n*–1)*k* to *nk*–1).

The block select circuits **106(1)–106(n)** control whether corresponding CAM blocks **102(1)–102(n)**, respectively, participate in compare operations. Each block select circuit **106** stores a class code for the corresponding CAM block **102** which may be used to selectively disable the CAM block from participating in, and therefore from affecting the results of, one or more compare operations. During a compare operation, a comparand word is provided to the CAM blocks **102** via CBUS, and a search code is provided to the

block select circuits **106(1)–106(n)** via bus SC. In alternate embodiments, the search code may be provided as part of the comparand word, in which case the CBUS is connected to the block select circuits **106**, or may be provided as part of a compare instruction. Each block select circuit **106** compares the received search code with its stored class code, and in response thereto, selectively disables the corresponding CAM block **102** from participating in the compare operation via a select signal SEL. In one embodiment, the block select circuit **106** enables its corresponding CAM block **102** to participate in the compare operation if the class code matches the search code and, conversely, disables the corresponding CAM block **102** if the class code does not match the search code. In alternate embodiments, more than one CAM block **102** may share the same block select circuit **106**.

The class codes assigned to the CAM blocks **102** may be used to partition the device **100** into individually selectable partitions of one or more CAM blocks **102**. For example, in one embodiment, data stored in the first CAM block **102(1)** may be assigned to a first class by storing a first class code in block select circuit **106(1)**, data stored in the second CAM block **102(2)** may be assigned to a second class by storing a second class code in block select circuit **106(2)**, and data stored in the remaining CAM blocks **102(3)–102(n)** may be assigned to a third class by storing a third class code in block select circuits **106(3)–106(n)**. Then, for example, data stored in the first CAM block **102(1)** may be selected for searching by setting the search code to match the first class code stored in the block select circuit **106(1)**.

When the search code matches the first class code, the block select circuit **106(1)** enables the first CAM block **102(1)** to compare the comparand word with its stored data corresponding to the first class code. If the search code does not match the second and third class codes, the remaining block select circuits **106(2)–106(n)** disable the corresponding, unselected CAM blocks **102(2)–102(n)**. When disabled, the unselected CAM blocks **102(2)–102(n)** do not drive the comparand word into their respective CAM arrays **104** for the compare operation, thereby precluding comparison with unselected data corresponding to the second and third class codes. In this manner, the CAM blocks **102(1)–102(n)** may be selectively searched according to class assignments, thereby allowing for a dynamically partitionable CAM device **100**.

Since the comparand word is not compared with data stored in the disabled CAM blocks **102(2)–102(n)**, the disabled CAM blocks **102(2)–102(n)** consume much less power during the compare operation than does the selected, enabled CAM block **102(1)**. In this manner, the class codes of present embodiments not only restrict compare operations to data in the selected CAM block(s), but also minimize power consumption of the unselected CAM block(s) during compare operations. The advantage of reduced power consumption in unselected CAM blocks during compare operations achieved by present embodiments may be particularly useful in applications where power consumption is a concern.

The ability to selectively enable or disable one or more CAM blocks from participating in compare operations may be especially useful for combining routing look-up functions for different classes of networks in a single device **100**. For example, in one embodiment, routing information for a first virtual private network (VPN) may be stored in a first CAM block **102(1)**, routing information for a second VPN may be stored in a second CAM block **102(2)**, routing information for a web search may be stored in a third block **102(3)**, and routing information for a local area network (LAN) may be



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stored in a fourth CAM block **102(4)**. Four unique class codes may be stored in corresponding block select circuits **106**. Of course, more than one CAM block may be assigned to a particular network by storing the appropriate class code in more than block select circuit **106**. During compare operations, comparand data corresponding to routing functions of one of these four networks may be exclusively compared with data stored in the corresponding CAM block(s) by simply setting the search code to match the appropriate class code. In some embodiments, an associative RAM may be partitioned into four partitions corresponding with the four class-defined partitions in the CAM device **100**.

FIG. 2 shows a CAM array **200** that is one embodiment of a CAM array **104** of FIG. 1. The array **200** includes a plurality of CAM cells **202** organized in any number of rows and columns. Each row of CAM cells **202** is coupled to a match line ML and a word line WL. Each word line WL is driven by an address decoder **204** to select one or more of CAM cells **202** for writing or reading. For alternative embodiments, multiple CAM blocks may share a decoder. Each match line ML provides the match results of a compare operation to the priority encoder **108** (see also FIG. 1). A match line ML indicates a match condition for the row only if all CAM cells **202** in that row match the comparand data. Each CAM cell **202** may be a binary, ternary, SRAM-based or DRAM-based CAM cell. In some embodiments, the match line ML is pre-charged for the compare operation. If any CAM cell **202** in the row does not match the comparand data, the CAM cell(s) **202** discharges the match line ML toward ground potential (e.g., logic low). Conversely, if all CAM cells **202** match the comparand data, the match line ML remains in a charged state (e.g., logic high). When the CAM block **102** is disabled in response to the select signal SEL, the comparand word is not driven into the array **200**, and the match lines ML may remain in their charged state during the compare operation, regardless if there is a mismatch. The match lines need not be pre-charged for a subsequent compare operation. The ability to maintain the match lines of unselected CAM blocks in their charged state during the compare operation may further reduce power consumption of present embodiments over prior art architectures.

Each column of CAM cells **202** is coupled to a bit line BL, a complementary bit line  $\overline{BL}$ , a comparand line CL, and a complementary comparand line  $\overline{CL}$ . The bit lines BL and  $\overline{BL}$  are coupled to sense amplifiers **206** that may enable data to be written to or read from a row of CAM cells **202**. The comparand lines CL and  $\overline{CL}$  are coupled to comparand drivers **208**, which in turn are coupled to a comparand register **210** via complementary data lines D and  $\overline{D}$ . The comparand drivers **208** selectively drive a comparand word received from the comparand register **210** via complementary data lines D and  $\overline{D}$  onto complementary comparand lines CL and  $\overline{CL}$  for comparison with data in CAM cells **202** in response the select signal SEL provided by the block select circuit **106**. The comparand register **210** may be shared by all CAM blocks **102(1)–102(n)**. As discussed above with respect to FIG. 1, the block select circuit **106** generates the select signal SEL in response to the search code and its stored class code.

In alternate embodiments, other CAM array architectures may be used. For example, in some embodiments, CAM array **200** may not include complementary comparand lines CL and  $\overline{CL}$ , in which case the complementary bit lines BL and  $\overline{BL}$  may be coupled to the comparand drivers **208** and be used to perform a compare operation as is generally

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known in the art. For example, in the first part of a compare cycle, compare data may be selectively driven onto BL and  $\overline{BL}$ , and during the second part of the compare cycle, BL and  $\overline{BL}$  may be driven with data to be output from CAM array **200**. For other embodiments, only one of comparand lines CL and  $\overline{CL}$  or bit lines BL and  $\overline{BL}$  may be needed.

FIG. 3 shows a block select circuit **300** that is one embodiment of the block select circuit **106**. The block select circuit **300** includes a memory **302** and a compare circuit **304**. The memory **302** stores the class code for the corresponding CAM block **102** of device **100**, and may be any suitable programmable memory element such as, for instance, a register, flip-flop, EEPROM, EPROM, SRAM, and so on. The compare circuit **304** compares the class code received from the memory **302** with a search code received from bus SC and, in response thereto, generates the select signal SEL which selectively enables or disables the corresponding CAM block **102**. The compare circuit **304** may be any suitable circuit which compares the search code and the class code, including for example an exclusive-OR type logic gate or a CAM cell.

FIG. 4 shows a block select circuit **400** that is one embodiment of the block select circuit **300**. The block select circuit **400** is shown to include a 3-bit memory **302** and a 3-bit compare circuit **304**, although in other embodiments more or less bits may be used. The memory **302** includes three data flip-flops **402(0)–402(2)**, and the compare circuit **304** includes three exclusive-NOR (XNOR) gates **404(0)–404(2)** and an AND gate **406**. Each XNOR gate **404(0)–404(2)** includes a first input terminal to receive a corresponding search code bit SC, a second input terminal to receive a corresponding class code bit CC from the corresponding flip-flop **402**, and an output terminal connected to the AND gate **406**. A 3-bit class code CC[0:2] may be clocked into respective flip-flops **402(0)–402(2)** using the clock signal CLK, where flip-flop **402(0)** stores the first class code bit CC[0], flip-flop **402(1)** stores the second class bit code CC[1], and flip-flop **402(2)** stores the third class code bit CC[2].

During compare operations, the XNOR gates **404(0)–404(2)** compare search code bits SC[0:2] with respective class code bits CC[0:2] and, if there is match, drive their output terminals to logic high. Conversely, if there is a mismatch, the XNOR gate **404** drives its output terminal to logic low. If all search code bits SC[0:2] match corresponding class code bits CC[0:2], then AND gate **406** asserts the select signal SEL to logic high, thereby enabling the corresponding CAM block **102** to participate in the compare operation. Otherwise, if any of the search code bits SC[0:2] mismatch corresponding class code bits CC[0:2], the AND gate **406** de-asserts the select signal to logic low, thereby disabling the corresponding CAM block **102** from participating in the compare operation. Since class code bits may be loaded into flip-flops **402(0)–402(2)** before a compare operation, the gate delay associated with generating the select signal during the compare operation is only 2 gate delays, one for XNOR gates **404** and one for AND gate **406**, and therefore has a negligible effect upon device performance.

FIG. 5 shows a block select circuit **500** that is another embodiment of the block select circuit **300**. Here, a logic circuit **502** is coupled to the output terminal of the AND gate **406** to allow for direct control of the select signal SEL using control signals EN and SEL\_OV. The logic circuit **502** includes an OR gate **504** having a first terminal coupled to the output terminal of the AND gate **406**, a second terminal to receive EN, and an output terminal coupled to a first input terminal of an AND gate **508**. The AND gate **508** includes

a second input terminal to receive SEL\_OV, and an output terminal to provide the select signal SEL. The signal EN enables the corresponding CAM block 102 to participate in the compare operation when the output of AND gate 406 is logic high. When asserted to logic high, EN enables the corresponding CAM block 102 for the compare operation and, conversely, when de-asserted to logic low, EN disables the corresponding CAM block 102 for the compare operation regardless if there is a match condition. The EN signal may be used to selectively disable CAM blocks 102, for instance, when defective. The signal EN is shown in FIG. 5 as being provided by a fuse 506, although in other embodiments EN may be provided by other means such as a programmable memory element, e.g., a register, flip-flop, EPROM, EEPROM, SRAM, etc. The signal SEL\_OV is a select override signal that, when asserted to logic high, may be used to force the select signal SEL to logic high to enable the corresponding CAM block 102 to participate in compare operations, irrespective of whether there is a class match. For an alternative embodiment, the relative locations of OR gate 504 and AND gate 508 may be reversed such that when SEL\_OV is set to a logic high state, then SEL will be set to a logic high state irrespective of whether there is a class match or the logic state of EN.

FIG. 6 shows a 1-bit comparand driver 600 that is used in one embodiment of the comparand drivers 208. Driver 600 includes AND gates 602, 604, and 606, and also includes buffers 608 and 610. AND gate 602 includes input terminals to receive the clock signal CLK and the select signal SEL, and an output terminal coupled to first input terminals of AND gates 604 and 606. AND gate 604 includes a second input terminal coupled to the data line D, and an output terminal coupled to the buffer 608, which in turn drives the comparand line CL. AND gate 606 includes a second input terminal coupled to the complementary data line  $\bar{D}$ , and an output terminal coupled to the buffer 610, which in turn drives the complementary comparand line  $\bar{CL}$ . Buffers 608 and 610 may be any suitable buffers to drive comparand data onto the comparand lines CL and  $\bar{CL}$ . A plurality of drivers 600 may share the AND gate 602.

During a compare operation, a comparand bit is provided to AND gate 604 via data line D, and a complementary comparand bit is provided to AND gate 606 via complementary data line  $\bar{D}$ . When CLK is logic high, the select signal SEL propagates through AND gate 602 to AND gates 604 and 606. If the select signal is asserted to logic high, AND gate 606 passes the comparand bit to the buffer 608, which in turn drives the comparand bit onto the comparand line CL. Similarly, AND gate 608 passes the complementary comparand bit to the buffer 610, which in turn drives the complementary comparand bit onto the complementary comparand line  $\bar{CL}$ . Thus, when the select signal SEL is asserted, the comparand driver 600 drives the comparand lines CL and  $\bar{CL}$  with the comparand data received from the comparand register 210 via data lines D and  $\bar{D}$ .

Conversely, if the select signal SEL is de-asserted to logic low to indicate that the corresponding CAM block 102 is not to participate in the compare operation, AND gates 606 and 608 force their respective output terminals to logic low. In response thereto, buffers 608 and 610 force the comparand line CL and the complementary comparand line  $\bar{CL}$ , respectively, to logic low. In this manner, when the select signal SEL is de-asserted, the comparand driver 600 does not drive complementary comparand data onto the comparand lines CL and  $\bar{CL}$ , thereby precluding the corresponding CAM block 102 from participating in the compare operation while minimizing power consumption in the CAM block.

The present invention is also particularly useful in increasing manufacturing yield of a CAM device by disabling defective CAM blocks in the device. Thus, for instance, during manufacture of a CAM device having n CAM blocks, if one or more of the CAM blocks are found to be defective or otherwise inoperable after manufacturing, rather than discarding the entire device, the defective blocks may be disabled using the block select circuits as described above, and the remaining non-defective CAM blocks may then be used for compare operations. For example, in one embodiment where the CAM device includes 8 CAM blocks each having 1k rows of CAM cells, if one of the CAM blocks is defective, that CAM block is disabled, and the remaining 7 CAM blocks may be used as a 7k CAM device. Accordingly, the ability to use the CAM device when one or more of its CAM blocks are defective advantageously increases manufacturing yield of the CAM device.

FIG. 7 shows a CAM device 700 that is a modified embodiment of the device 100 of FIG. 1 which allows for one or more defective CAM blocks to be disabled for CAM operations, and also includes circuitry which translates or re-assigns address locations in defective CAM blocks to address locations in non-defective CAM blocks. The device 700 includes address logic 701, a plurality of CAM blocks 702(1)–702(n), a plurality of block select circuits 706(1)–706(n) corresponding to CAM blocks 702(1)–702(n), respectively, a priority encoder 708, and full flag logic 710. Each of the block select circuits 706(1)–706(n) provides to the corresponding CAM block 702 a select signal which may be used as described above to disable the CAM block 702 if, for example, the CAM block 702 is defective.

The block select circuit 706 may be any suitable circuit to provide either a logic high (enabling) or a logic low (disabling) select signal SEL to the corresponding CAM block 702. In some embodiments, the block select circuit 706 includes a memory (not shown in FIG. 7) for storing a binary value indicative of SEL. In some embodiments, the block select circuit 706 provides a logic high SEL signal if the corresponding CAM block 702 is not defective, and provides a logic low SEL signal if the corresponding CAM block 702 is defective. In one embodiment, the block select circuit 706 may include the block select circuit 500 (FIG. 5), in which case the signal EN may be set to a low logic state by blowing fuse 506 to disable a defective CAM block 702 via signal SEL. In other embodiments, the block select circuit 706 may be a fuse (or a memory element) connected between the CAM block 702 and a voltage supply, in which case the fuse may be blown to provide a logic low SEL signal to disable the corresponding CAM block 702.

After fabricating the device 700, each of the CAM blocks 702(1)–702(n) is tested in a suitable manner. For each CAM block 702 that is found to be defective, the corresponding block select circuit 706 is configured to provide a logic low select signal SEL to the CAM block 702 to disable the CAM block. Conversely, for each CAM block 702 that is not defective, the corresponding block select circuit 706 is configured to provide a logic high select signal to the CAM block to enable its participation in CAM operations.

Testing the CAM blocks of a CAM device and then selectively disabling the defective CAM blocks in an embodiment using a fuse to provide SEL is illustrated with reference to the flow chart of FIG. 8. Here, a fuse (not shown for simplicity) in each block select circuit 706 is coupled to a voltage supply and thus initially provides an asserted (e.g., logic high) SEL to enable the corresponding CAM block 702. Each CAM block 702 is tested in a suitable manner to determine whether it is defective (step 750). If the CAM

block is defective, as tested at step 751, SEL is de-asserted (e.g., to logic low) to disable the defective CAM block by blowing the fuse. Otherwise, if the CAM block is not defective, the corresponding fuse is not blown, and the corresponding CAM block remains enabled. If all CAM blocks have been tested, as determined at step 753, processing is finished (step 754). Otherwise, the next CAM block is tested and thereafter disabled if found to be defective (steps 750–752).

During a compare operation, each CAM block 702 receives comparand data from the comparand bus CBUS in a manner similar to that of CAM blocks 102 of device 100 of FIG. 1. Other signals provided to device 700 during the compare operation may be a clock signal CLK, one or more instructions from an instruction decoder (not shown for simplicity), and other control signals. Each CAM block 702 provides a plurality of match line signals to the priority encoder 708 via corresponding match lines ML. The match lines carry match signals indicative of match conditions in the CAM arrays 704. For simplicity, the plurality of match lines ML from each CAM block 702 are represented collectively in FIG. 7. The priority encoder 708 generates an index corresponding to one of the matching CAM words in the device 700, which as described above may be index of the highest-priority matching CAM row.

Each CAM block 702 provides a full flag signal FF indicative of whether the CAM block is full, i.e., whether there are any available row in the CAM block 702 to store data, to full flag logic 710. The full flag signal FF may be generated for each CAM block 702 in a well-known manner using one or more valid bits in each row of the CAM block. The full flag signals FF<sub>1</sub> to FF<sub>n</sub> provided by CAM blocks 702(1)–702(n), respectively, are combined in a well-known manner in full flag logic 710 to generate a device full flag, FF<sub>device</sub>, indicative of whether there are any available rows in the device 700. When a CAM block 702 is found to be defective or otherwise inoperable for its intended purpose, the CAM block 702 is configured to maintain an asserted full flag signal FF to indicate that the defective CAM block 702 does not include any available memory locations. In one embodiment, the full flag signal FF for the defective CAM block may be maintained in the asserted state by forcing the valid bits in its array 704 to an asserted state. In other embodiments, a fuse may be provided within or associated with each CAM block 702 that, when blown, forces the corresponding full signal FF to be asserted.

Address logic 701 is shown in FIG. 7 as coupled to an address bus ABUS and each of the CAM blocks 702(1)–702(n). During read and write operations, an address provided to the device 700 may be received into address logic 701 via address bus ABUS, and thereafter used to select a row in one of the CAM blocks 702(1)–702(n) for the read or write operation. In accordance with the present invention, if a CAM block 702 to which the address refers is defective, and is thus disabled for the operation using the corresponding block select circuit 706 as describe above, address logic 701 translates the address from the defective or disabled CAM block to a non-defective CAM block. Conversely, if the CAM block 702 to which the address refers is non-defective, and is thus enabled for operation, address logic 701 forwards the address to the appropriate CAM block 702. As explained more fully below, address logic 701 ensures a contiguous addressing scheme in the CAM blocks 702 when one or more CAM blocks 702 are defective and disabled, even when the non-defective CAM block(s) 702 are not adjacent to each other.

For alternate embodiments, address logic 701 may be omitted. For one example, contiguous non-defective blocks starting from block 702(1) may still be used. For other embodiments, any non-defective block may be used.

FIG. 9 shows a CAM device 800 that is one embodiment of the CAM device 700. CAM device 800 is shown to include address translation logic 801, four CAM blocks 802(0)–802(3), and a main priority encoder 806. Each CAM block 802 includes a CAM array 704 (e.g., a 1k CAM array), a block priority encoder 804, and match flag logic 805. Of course, in other embodiments, there may be any number of CAM blocks 802, and each CAM block array 704 may include any number of rows of CAM cells. The address A may include any suitable number of bits. In the embodiment of FIG. 9, the function of the priority encoder 708 of FIG. 7 is distributed between the individual block priority encoders 804 within the CAM blocks 802(0)–802(3) and the main priority encoder 806. During a read or write operation, a 14-bit address A[13:0] may be provided to the device 800 via the address bus ABUS. The first two address bits A[13:12] are the block address bits and are provided to address translation logic 801, which in turn selects one of the CAM blocks 802(0)–802(3) for the read or write operation via block select signals BS<sub>0</sub> to BS<sub>3</sub>, respectively. The remaining twelve address bits, A[11:0], select a row in the CAM array 704 selected by address translation logic 801 for the operation, and may be provided to each CAM block 802. During read or write operations, data may be read from or written to the row identified by row address bits A[11:0] in the CAM block 802 selected by address translation logic 801.

Information indicative of which CAM blocks 802 are found to be defective during testing may be used to configure address translation logic 801 to re-address the non-defective CAM blocks 802 so as to occupy, for instance, the contiguous highest-priority address space (e.g., the lowest numbered addresses). During a read or write operation, address translation logic 801 receives block address bits A[13:12]. If a CAM block 802 selected by block address bits A[13:12] is non-defective or otherwise enabled, address translation logic 801 asserts the corresponding block select signal BS to enable the selected CAM block 802 for the operation. For example, if an address [13:0] selects the first row in the first CAM block 802(0) for reading, and CAM block 802(0) is non-defective, address translation logic 801 asserts BS<sub>0</sub> to logic high while maintaining BS<sub>1</sub>, BS<sub>2</sub>, and BS<sub>3</sub> in a logic low, de-asserted state. The asserted BS<sub>0</sub> signal causes row address bits A[11:0] to be latched into the first CAM block 802(0), thereby facilitating a read from the first CAM block 802(0).

Conversely, if a CAM block 802 selected by block address bits A[13:12] is defective or otherwise disabled, address translation logic 801 selects another CAM block for the operation by asserting its corresponding block select signal BS. For example, if the address [13:0] selects the first row in the first CAM block 802(0) for reading, and CAM block 802(0) is defective and the second CAM block 802(1) is non-defective, address translation logic 801 may assert BS<sub>1</sub> to logic high while maintaining BS<sub>0</sub>, BS<sub>2</sub>, and BS<sub>3</sub> in a logic low, de-asserted state. The asserted BS<sub>1</sub> signal causes row address bits A[11:0] to be latched into the second CAM block 802(1), thereby facilitating a read from the second CAM block 802(1). In this manner, address translation logic 801 may re-address read or write operations from defective CAM blocks to non-defective CAM blocks.

In some embodiments, the block select signal BS provided to the CAM block 802 may be used as an address

gating signal to facilitate address translation in accordance with present embodiments. For example, FIG. 10 shows a CAM array 900 that is one embodiment of the array 704 of FIG. 7. The array 900 includes a plurality of CAM cells 202 organized in any number of rows and columns, and operates in a manner similar to the CAM array 200 described above with respect to FIG. 2. That is, during compare operations, comparand data provided by the comparand register 210 is selectively driven onto the complementary comparand lines CL and  $\overline{\text{CL}}$  in response to the select signal SEL provided by the block select circuit 706. If the array 900 is non-defective, the select signal SEL is asserted to logic high to allow the comparand word to be driven into the array 900 for comparison with CAM words stored therein. Conversely, if the array 900 is found to be defective during testing, the block select circuit 706 is configured to provide a de-asserted select signal SEL to the comparand drivers 208 to prevent comparand data from being driven onto the comparand lines CL and  $\overline{\text{CL}}$ , thereby disabling the array 900.

Address bits A[11:0] are provided from address bus ABUS to the address decoder 204. Address gating logic 902 is connected between the address decoder 204 and corresponding word lines WL of the array 900 via gated lines GL, and selectively drives a word line WL identified by A[11:0] in response to the block select signal BS. For example, during a read or write operation, address decoder 204 decodes A[11:0] to select a row of CAM cells 202 for the operation, and drives a corresponding gated line GL to logic high. If BS is asserted to logic high, address gating logic 902 drives the corresponding word line WL to select the row of CAM cells 202 for the operation. Conversely, if BS is de-asserted to logic low, address logic 902 does not drive any of the word lines WL to logic high, regardless of A[11:0], thereby preventing CAM cells 202 in the array from being addressed for the operation.

In one embodiment, address gating logic 902 may include for each word line WL in the array 900 an AND gate (not shown) having an output terminal coupled to the word line, a first input terminal coupled to the corresponding gated line GL, and a second input terminal to receive the block select signal BS. In this manner, the AND gates may be used to selectively gate the addressing of CAM cells in the block in response to BS. Of course, in other embodiments other suitable logic may be used.

FIG. 11 shows address translation logic 1000 that is one embodiment of the address translation logic 801 of FIG. 9. Logic 1000 includes decode logic 1002, four 4-input multiplexers 1004(0)–1004(3), and four corresponding memory elements 1006(0)–1006(3), respectively. Decode logic 1002 has an input terminal to receive block address bits A[13:12], and has four output terminals coupled to corresponding input terminals of the multiplexers 1004(0)–1004(3) via lines 1008(0)–1008(3), respectively. Decode logic 1002 decodes block address bits A[13:12], and in response thereto, asserts one of output lines 1008(0)–1008(3) to logic high. For example, if address bits A[13:12] are “00”, which is equivalent to the decimal value “0”, decode logic asserts line 1008(0); if address bits A[13:12] are “01”, which is equivalent to the decimal value “1”, decode logic 1002 asserts line 1008(1); if address bits A[13:12] are “10”, which is equivalent to the decimal value “2”, decode logic 1002 asserts line 1008(2); and if address bits A[13:12] are “11”, which is equivalent to the decimal value “3”, decode logic 1002 asserts line 1008(3).

The multiplexers 1004(0)–1004(3) each include an output terminal coupled to a corresponding one of the CAM blocks 802(0)–802(3), respectively, and a control terminal coupled

to a corresponding one of the memory elements 1006(0)–1006(3), respectively. Each memory element 1006 stores address translation information that when provided to the corresponding multiplexer 1004 selects one of the signals provided by decode logic 1002 to be output as the block select signal BS. In this manner, multiplexers 1004(0)–1004(3) may dynamically assign block address values to CAM blocks 802(0)–802(3), respectively.

In accordance with present embodiments, a read or write operation to a defective CAM block may be re-addressed to a non-defective CAM block by manipulating the address translation information stored in the memory elements 1006(0)–1006(3). In some embodiments, where all CAM blocks 802(0)–802(3) are non-defective, the memory elements 1006(0)–1006(3) store default block address values so as to not alter CAM addressing during the read or write operation. That is, the default block address values cause respective multiplexers 1004(0)–1004(3) to select corresponding signals on lines 1008(0)–1008(3) as block select signals BS\_0 to BS\_3, respectively. For example, memory element 1006(0) may store a default block address value of “0” to cause multiplexer 1004(0) to select the signal on line 1008(0) as BS\_0, memory element 1006(1) may store a default block address value of “1” to cause multiplexer 1004(1) to select the signal on line 1008(1) as BS\_1, memory element 1006(2) may store a default block address value of “2” to cause multiplexer 1004(2) to select the signal on line 1008(2) as BS\_2, and memory element 1006(3) may store a default block address value of “3” to cause multiplexer 1004(3) to select the signal on line 1008(3) as BS\_3. In this manner, address translation logic 1000 selects for the read or write operation the CAM block identified by address bits A[13:12]. Table 1 summarizes the four default block address/multiplexer select values (MUX) and corresponding address space when all CAM blocks are non-defective.

TABLE 1

Block	Status	MUX	address space
CAM 802(0)	non-defective	0	0 to k-1
CAM 802(1)	non-defective	1	k to 2k-1
CAM 802(2)	non-defective	2	2k to 3k-1
CAM 802(3)	non-defective	3	3k to 4k-1

The CAM device 800 is then tested to determine if any CAM blocks are defective. Where it is determined that one or more CAM blocks are defective, the select values stored in memory elements 1006(0)–1006(3) may be modified to re-address the non-defective CAM blocks. For example, if after testing it is determined that CAM block 802(0) is defective, and is thereafter disabled using the corresponding block select circuit 706 as described above (see also FIG. 7), the 1k CAM rows in the defective CAM block 802(0) are no longer available, and therefore the device 800 now has only 3k available CAM rows available, i.e., 1k rows in each of the 3 non-defective CAM blocks 802(1)–802(3). Since the first CAM block 802(0) is not available, it is desirable for the second CAM block 802(1) to be the highest-priority CAM block (e.g., having address space 0 to k-1), for the third CAM block 802(2) to be the second highest-priority CAM block (e.g., having address space k to 2k-1), and for the fourth CAM block 802(3) to be the third highest-priority CAM block (e.g., having address space 2k to 3k-1).

The block address values stored in corresponding memory elements 1006(0)–1006(3) may be modified to implement a new addressing scheme for the non-defective CAM blocks 802(1)–802(3). For example, in one embodiment, the block

address value stored in the memory element **1006(1)** is set to “0” so that multiplexer **1004(1)** selects the signal on line **1008(0)** as BS<sub>1</sub> to be provided to CAM block **802(1)**. When A[13:12] equals “00”, decode logic **1002** asserts line **1008(0)** to logic high, which in turn now passes through multiplexer **1006(1)** to select the second CAM block **802(1)** for the operation. In this manner, address translation logic **1000** translates address space 0 to k-1 from CAM block **802(0)** to CAM block **802(1)**.

Similarly, the block address value stored in the memory element **1006(2)** is set to “1” so that multiplexer **1004(2)** selects the signal on line **1008(1)** as BS<sub>2</sub> to provide to CAM block **802(2)**. When A[13:12] equals “01”, decode logic **1002** asserts line **1008(1)** to logic high, which in turn now passes through multiplexer **1004(2)** to select the third CAM block **802(2)** for the operation, thereby translating address space k to 2k-1 from CAM block **802(1)** to CAM block **802(2)**. Similarly, the block address value stored in the memory element **1006(3)** is set to “2” so that multiplexer **1004(3)** selects the signal on line **1008(2)** as BS<sub>3</sub> to provide to CAM block **802(3)**. When A[13:12] equals “10”, decode logic **1002** asserts line **1008(2)** to logic high, which in turn now passes through multiplexer **1004(3)** to select the fourth CAM block **802(3)** for the operation, thereby translating address space 2k to 3k-1 from CAM block **802(2)** to CAM block **802(3)**. Table 2 summarizes the select values when CAM block **802(0)** is defective and CAM blocks **802(1)–802(3)** are non-defective.

TABLE 2

Block	Status	MUX	address space
CAM 802(0)	defective	3	3k to 4k-1*
CAM 802(1)	non-defective	0	0 to k-1
CAM 802(2)	non-defective	1	k to 2k-1
CAM 802(3)	non-defective	2	2k to 3k-1

\*not used

By translating address space in CAM blocks **802(1)–802(3)**, respectively, present embodiments may re-address rows in non-defective CAM blocks **802(1)–802(3)** with the highest-priority CAM addresses, e.g., row addresses 0 to 3k-1. In this manner, the three non-defective CAM blocks **802(1)–802(3)** of device **800** may be sold and operated as a 3k CAM array. This is in contrast to prior art CAM devices, which are typically discarded if any of the CAM blocks therein are found to be defective. The ability to re-address the defective CAM block **802(0)** and use the non-defective CAM blocks **802(1)–802(3)** of device **800** as a 3k CAM array, rather than discarding the device **800**, may significantly increase manufacturing yield.

In the example above, address space in the defective CAM block **802(0)** is translated from row assignments 0 to k-1 to row assignments 3k to 4k-1 by changing the select value stored in the memory element **1006(0)** from “0” to “3”. This ensures that the defective CAM block **802(0)** will not be addressed during read or write operations. That is, since the 3 non-defective CAM blocks are used as a 3k CAM array having address space 0 to 3k-1, address space higher than 3k-1 is not used, and therefore the defective CAM block **802(0)** will not be addressed. For an alternate embodiment, the block select signal BS<sub>0</sub> can be set to a low logic state to disable block **802(0)**. For one example, the output of each MUX can be coupled to a logic circuit (e.g., one or more AND, OR, XOR, NOT circuits) and memory **1006** configured to disable BS and its corresponding block when a particular value is programmed into memory **1006** (or the value is changed in memory **1006**).

In other embodiments, the block address values stored in memory elements **1006(0)–1006(3)** may be modified to translate address space in any number n of non-defective CAM blocks into a contiguous address space of 0 to (k)n-1, irrespective of whether the non-defective CAM blocks are adjacent to one another. Thus, for example, if in one embodiment the CAM blocks **802(0)** and **802(2)** are defective and CAM blocks **802(1)** and **802(3)** are non-defective, the non-defective CAM blocks **802(1)** and **802(3)** may be configured for operation as a 2k CAM array by setting the block address values for memory elements **1006(1)** and **1006(3)** to “0” and “1”, respectively. In this manner, the first 1k address space corresponding A[13:12] equal to “00” selects CAM block **802(1)**, and the second 1k address space corresponding to A[13:12] equal to “01” selects CAM block **802(3)**. The block address values stored in memory elements **1006(0)** and **1006(2)** each may be either “2” or “3” to preclude their selection during operation, since addresses above 2k-1, i.e., the third or fourth 1k address spaces corresponding to A[13:12] equal to “10” or “11”, respectively, are not used. Table 3 summarizes the block address values and corresponding address space when CAM blocks **802(0)** and **802(2)** are defective and CAM blocks **802(1)** and **802(3)** are non-defective.

TABLE 3

Block	Status	MUX	address space
CAM 802(0)	defective	2 or 3	>2k*
CAM 802(1)	non-defective	0	0 to k-1
CAM 802(2)	defective	2 or 3	>2k*
CAM 802(3)	non-defective	1	k to 2k-1

\*not used

Each memory element **1006** may be any suitable structure to provide a block address value to the corresponding multiplexer **1004** to select one of lines **1008(0)–1008(3)** to pass as the block select signal BS. In some embodiments, the memory element may be a flip-flop, register, look-up table, or non-volatile memory such as EPROM or Flash memory. In other embodiments, the memory element **1006** may include one or more fuses to provide the block address value to the corresponding multiplexer **1004**.

For one example, in one embodiment of the CAM device **800**, each memory element **1006** includes two fuses coupled to a voltage supply to initially provide the binary value “11” to corresponding multiplexers **1004**. In this example, since each multiplexer **1004** initially selects the signal line connected to its “3” input in response to the binary block address value “11”, the “3” input of each multiplexer **1004** is connected to a corresponding numbered signal line **1008** from decode logic **1002**. That is, input 3 of multiplexer **1004(0)** is connected to the line **1008(0)**, input number 3 of multiplexer **1004(1)** is connected to the line **1008(1)**, input number 3 of multiplexer **1004(2)** is connected to the line **1008(2)**, and input number 3 of multiplexer **1004(3)** is connected to the line **1008(3)**. The remaining multiplexer inputs 0, 1, and 2 may be connected to lines **1008** in any suitable configuration. In this manner, multiplexer **1004(0)** asserts BS<sub>0</sub> when A[13:12] equals “00”, multiplexer **1004(1)** asserts BS<sub>1</sub> when A[13:12] equals “01”, multiplexer **1004(2)** asserts BS<sub>2</sub> when A[13:12] equals “10”, and multiplexer **1004(3)** asserts BS<sub>3</sub> when A[13:12] equals “11”. Then, if after testing one or more of the CAM blocks **802** are found to be defective, the fuses of each memory element **1006** may be selectively blown to translate address space from defective CAM blocks to non-defective CAM

blocks **802** to facilitate contiguous addressing in the manner described above. Additionally, the two logic ones may be ANDed together and provided as one input to an AND gate, and the other input to the AND gate coupled to a respective BS signal output by each MUX. When a defective row is programmed to a value other than "11", then the respective BS signal and corresponding block will be disabled.

In order to maintain address consistency between read or write operations and compare operations when address space in a defective CAM block is translated to a non-defective CAM block, address translation information used during the read or write operation is also used to calculate the address or index of a matching CAM row during compare operations. Thus, for example, if address space 0 to k-1 is translated from CAM block **802(0)** to **802(1)**, and there is a match in CAM block **802(1)** during a subsequent compare operation, the priority encoder **806** ensures that the matching index from CAM block **802(1)** lies within address space 0 to k-1, rather than within address space k to 2k-1. In this manner, address translations facilitated during a read or write operation are reflected during subsequent compare operations.

Referring again to FIG. 9, during compare operations, a comparand word provided on CBUS is compared to data stored in all enabled (e.g., non-defective) CAM blocks **802**. For each enabled CAM block **802**, if there is a match condition in response to the compare operation, match flag logic **805** asserts a match flag (MF\_0 to MF\_3) to a logic high state, and the priority encoder **804** within the CAM block **802** outputs the 12-bit row index I (or address) of the highest priority matching CAM row in the block. If there is not a match, the match flags are not asserted (i.e., match flag logic **805** sets MF to low logic state). For one embodiment, each match flag logic includes a programmable element (e.g., a fuse or other memory element) that is programmed when the CAM block is disabled. The match flags MF\_0 to MF\_3 and row indexes I0-I3 from CAM blocks **802(0)**-**802(3)**, respectively, are provided to the main priority encoder **806**. The main priority encoder **806** adds a unique block index to each row index I provided by CAM blocks **802(0)**-**802(3)** to form a corresponding device index. The main priority encoder **806** uses the match flag signals MF\_0 to MF\_3 to select the highest-priority device index from CAM blocks **802(0)**-**802(3)** to output as the system index, I\_sys.

The main priority encoder **806** is programmable and stores the block indexes for CAM blocks **802(0)**-**802(3)** in memory (not shown in FIG. 9). The block indexes are dynamic values that may be modified or programmed to reflect and thus maintain consistency with address translations as described above in a read or write operation. Initially, the main priority encoder stores a block index of "00" for CAM block **802(0)**, a block index of "01" for CAM block **802(1)**, a block index of "10" for CAM block **802(2)**, and a block index of "11" for CAM block **802(3)**. These initial block indexes, which are used when all CAM blocks **802(0)**-**802(3)** are enabled, mirror the block address values stored in memory elements **1006(0)**-**1006(3)** of address translation logic **1000** of FIG. 11. If one or more CAM blocks **802** are found to be defective or are otherwise disabled, the block indexes stored in main priority encoder **806** are modified to reflect address translations during the read or write operation. For example, if CAM blocks **802(0)** and **802(2)** are disabled and address spaces in non-defective CAM blocks **802(1)** and **802(3)** are translated to address space 0 to k-1 and address space k to 2k-1, respectively, main priority encoder **806** adds a block index of "00" to row

index **I1** to generate the device index for CAM block **802(1)** and adds a block index "01" to row index **I3** to generate the device index for CAM block **802(3)**. In this manner, address consistency between read/write operations and compare operations is maintained.

FIG. 12 shows a priority encoder **1200** that is one embodiment of the main priority encoder **806** of FIG. 9. The priority encoder **1200** includes a chain of four multiplexers **1202(0)**-**1202(3)**, four corresponding memory elements **1204(0)**-**1204(3)**, and select logic including an inverter **1206** and OR gates **1208** and **1210**. Each memory element **1204** stores a 2-bit block index for a corresponding CAM block **802**. The memory elements **1204** may be the same as memory elements **1006** of FIG. 11, or they may be separate memory elements. Each multiplexer **1202** includes a first input (i.e., the "1" input) coupled to the output of a preceding multiplexer **1202** in the chain, a second input (i.e., the "0" input) to receive a concatenation of a 12-bit index I of the highest priority match (if any) from a corresponding CAM block **802** and a 2-bit block index from the corresponding memory element **1204**, an output coupled to input "1" of a next multiplexer **1202**, and a select terminal to receive match information from the CAM blocks **802**.

The concatenation of a 12-bit row index I and the 2-bit block index forms a 14-bit device index of the highest-priority match, if any, from a corresponding CAM block **802**. The 1 input of the first multiplexer **1202(0)** receives a default binary "0" value. MF\_0 is inverted by inverter **1206** and provided as the select signal to multiplexer **1202(0)**, and provided directly as the select signal to multiplexer **1202(1)**. MF\_0 and MF\_1 are combined in OR gate **1208** and provided as the select signal for multiplexer **1202(2)**. MF\_2 and the result from OR gate **1208** (i.e., MF\_0 + MF\_1, where + is the logic OR function) are combined in OR gate **1210** and provided as the select signal for multiplexer **1202(3)**. As explained below, the match flags MF control whether each multiplexer **1202** passes a concatenated device index from a previous CAM block or the concatenated device index of the corresponding CAM block.

In this example, CAM block **802(0)** is the highest-priority block, CAM block **802(1)** is the next highest-priority block, and so on. For each multiplexer stage, if there is a match in the corresponding CAM block **802**, the row index I and block index are forwarded to the next stage if there is not a match condition in a previous or higher-priority CAM block **802**. If there is a match condition in a higher-priority CAM block **802**, the row index I plus block index from the higher-priority CAM block are forwarded to the next stage.

For example, if there is a match condition in the first CAM block **802(0)**, priority encoder **804** of CAM block **802(0)** provides the 12-bit row index I0 of its highest-priority match to input 0 of multiplexer **1202(0)**, where it is concatenated with the block index from memory element **1204(0)** to generate the device index for CAM block **802(0)**. The match flag MF\_0 is asserted to logic high to indicate the match condition. In response thereto, inverter **106** provides a logic low or "0" select signal to multiplexer **1202(0)**, which in turn forwards the device index from CAM block **802(0)** to the next multiplexer **1204(1)**. The logic high MF\_0 signal causes multiplexer **1204(1)** to select input 1, and thus forwards the device index from CAM block **802(0)** to the next multiplexer **1202(2)**. MF\_0 ripples through OR gates **1208** and **1210** and causes multiplexers **1202(2)** and **1202(3)** to output the device index from CAM block **802(0)** as I\_sys.

Maintaining equivalent values in corresponding memory elements **1006(0)**-**1006(3)** and **1204(0)**-**1204(3)** ensures

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addressing consistency between read or write operations and compare operations. For instance, in one embodiment where all CAM blocks **802(0)–802(3)** are non-defective or otherwise enabled, memory elements **1204(0)–1204(3)** store values of “00”, “01”, “10” and “11”, respectively. In this manner, “00” is added to row index **10** so that the first 1k addresses are mapped to the first block **802(0)**, “01” is added to row index **11** so that the second 1k addresses are mapped to the second block **802(1)**, “10” is added to row index **12** so that the third 1k addresses are mapped to the third block **802(2)**, and “11” is added to row index **13** from block **802(3)** so that the fourth 1k addresses are mapped to the fourth block **802(3)**.

The block indexes stored in memory elements **1204(0)–1204(3)** may be changed when address space in one or more CAM blocks **802** is translated to maintain addressing consistency. For example, in one embodiment where CAM block **802(0)** is defective, CAM blocks **802(1)–802(3)** may be configured to operate as a 3k CAM array as described above with respect to address translation logic **1000** (FIG. **11**) by setting block address values or their binary equivalents of “0”, “1” and “2” into memory elements **1006(1)**, **1006(2)**, and **1006(3)**, respectively. This configures CAM block **802(1)** to have the highest-priority address space, i.e., addresses 0 to k–1, CAM block **802(2)** to have the next highest-priority address space, i.e., k to 2k–1, and CAM block **802(3)** to have the lowest-priority address space, i.e., addresses 2k to 2k–1.

In accordance with present embodiments, the block indexes stored in memory elements **1204(0)–1204(3)** are modified to reflect address translations facilitated in address translation logic **1000**. Since CAM block **802(0)** is disabled (and thus does not require address space), the block indexes of the remaining enabled CAM blocks **802(1)–802(3)** may be modified to re-assign block priority in the CAM device **800**. For example, the block index stored in memory element **1204(1)** may be set to “00” so that when concatenated with row index **11** from the highest-priority CAM block **802(1)**, the resultant device index corresponds to the highest-priority address space, i.e., addresses 0 to k–1. Similarly, the block indexes stored in memory elements **1204(2)** and **1204(3)** may be modified to “01” and “10”, respectively, to reflect address spaces k to 2k–1 and 2k to 3k–1, respectively. Since in this embodiment addresses larger than 3k are not used, the block index stored in memory element **1204(0)**, which corresponds to the defective CAM block **802(0)**, may be set to “11” so that disabled CAM block **802(0)** is not addressed. Note that the match flag signal for a disabled CAM block will be set to a low logic state.

I claim:

1. A content addressable memory (CAM) device comprising:

a plurality of CAM blocks each having an array of CAM cells organized in a number of rows and columns; and select means for selectively disabling one or more of the plurality of CAM blocks during a compare operation.

2. The device of claim 1, wherein the CAM blocks comprise binary CAMs.

3. The device claim 1, wherein the CAM blocks comprise ternary CAMs.

4. The device claim 1, further comprising a priority encoder coupled to each of the plurality of CAM blocks.

5. The device of claim 1, wherein the select means comprises a plurality of block select circuits, each of which corresponds to and selectively disables a corresponding CAM block during the compare operation.

6. The device of claim 5, wherein each block select circuit comprises:

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a memory for storing a class code for the corresponding CAM block; and

a compare circuit for comparing a search code with the class code to generate a select signal, the select signal for selectively disabling the corresponding CAM block during the compare operation.

7. The device of claim 6, wherein the memory comprises a register.

8. The device of claim 6, wherein the compare circuit comprises one or more exclusive-OR type logic gates.

9. The device of claim 1, wherein each CAM block further comprises:

a number of comparand drivers each coupled to a corresponding column of the array and coupled to receive comparand data.

10. The device of claim 9, wherein each CAM block further comprises a number of complementary comparand lines corresponding to the number of columns, the comparand drivers coupled to the complementary comparand lines.

11. A content addressable memory (CAM) device comprising:

a plurality of CAM blocks each including an array of CAM cells; and

a plurality of block select circuits, each comprising:

a memory for storing a class code for a corresponding CAM block; and

a compare circuit having a first input to receive a search code, a second input to receive the class code from the memory, and an output coupled to the corresponding CAM block.

12. The device of claim 11, wherein each block select circuit for selectively disabling a corresponding CAM block during a compare operation in response to a comparison between the search code and the class code.

13. The device of claim 11, wherein each CAM block further comprises:

a number of comparand line drivers, each coupled to a corresponding column of the array and to a corresponding block select circuit, each block select circuit for selectively disabling one or more comparand line drivers of the corresponding CAM block.

14. The device of claim 11, wherein each CAM block further comprises:

an array of CAM cells arranged in a plurality of rows and columns;

a plurality of lines, each coupled to the CAM cells in a corresponding column of the array;

a plurality of inputs to receive a comparand word from a comparand bus; and

a plurality of comparand drivers, each coupled to a corresponding one of the plurality of lines and having an input terminal coupled to the block select circuit, wherein the plurality of comparand drivers for selectively driving a comparand word onto the lines in response to the select signal provided by the block select circuit.

15. The device of claim 14, wherein the lines comprise comparand lines.

16. A method of selectively disabling one or more of a plurality of content addressable memory (CAM) blocks of a CAM device from participating in a compare operation between a comparand word and data stored in respective arrays of the CAM blocks, the method comprising:

assigning a class code to each CAM block;

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providing a search code to the CAM blocks;  
for each CAM block, comparing the search code with the  
class code to generate a select signal; and  
selectively disabling one or more of the CAM blocks in  
response to the select signal.

17. The method of claim 16, wherein the class codes  
define individually searchable partitions in the CAM device.

18. The method of claim 16, wherein a unique class code  
is assigned to each CAM block.

19. The method of claim 16, wherein a same class code is  
assigned to more than one CAM block.

20. The method of claim 16, wherein the CAM block is  
disabled during the compare operation if the search code  
does not match the class code.

21. The method of claim 16, wherein the CAM block is  
not disabled during the compare operation if the search code  
matches the class code.

22. The method of claim 16, wherein selectively disabling  
the CAM block comprises preventing the comparand word  
from being driven into the array.

23. The method of claim 22, wherein preventing the  
comparand word from being driven into its array comprises  
forcing a set of comparand lines to a predetermined state.

24. A content addressable memory (CAM) device com-  
prising:

a plurality of CAM blocks each having an array of CAM  
cells organized in a number of rows and columns, and  
comprising a number of complementary comparand  
lines corresponding to the number of columns;

select means for selectively disabling one or more of the  
plurality of CAM blocks during a compare operation;  
and

a number of comparand drivers each coupled to a corre-  
sponding complementary comparand line and coupled to  
receive comparand data.

25. The device of claim 24, wherein the select means  
comprises a plurality of block select circuits, each of which  
corresponds to and selectively disables a corresponding  
CAM block during the compare operation.

26. The device of claim 25, wherein each block select  
circuit comprises:

a memory for storing a class code for the corresponding  
CAM block; and

a compare circuit for comparing a search code with the  
class code to generate a select signal, the select signal  
for selectively disabling the corresponding CAM block  
during the compare operation.

27. The device of claim 26, wherein the memory com-  
prises a register.

28. A content addressable memory (CAM) device com-  
prising:

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a plurality of CAM blocks each having an array of CAM  
cells organized in a number of rows and columns;

a plurality of comparand drivers each for providing com-  
parand data to a corresponding CAM block;

a plurality of memory circuits each for storing a class code  
indicative of a class of data for storage in a correspond-  
ing CAM block; and

a plurality of compare circuits each for comparing a  
search key with one of the class codes for a corre-  
sponding CAM block, and operable to enable the  
plurality of comparand drivers to provide the com-  
parand data to the corresponding CAM block in  
response to the comparison.

29. The device of claim 28, wherein the compare circuits  
are further operable to disable the plurality of comparand  
drivers from providing the comparand data to the corre-  
sponding CAM block in response to the comparison.

30. The device of claim 29, wherein the disabled com-  
parand drivers force the comparand data to a predetermined  
state.

31. A content addressable memory (CAM) device com-  
prising:

a first CAM block having an array of CAM cells orga-  
nized in a number of rows and columns, and for storing  
data of a first class type;

a second CAM block having an array of CAM cells  
organized in a number of rows and columns, and for  
storing data of a second class type; and

means for enabling comparand data for comparison with  
the data of the first class type of the first CAM block  
and for disabling the comparand data for comparison  
with the data of the second class type of the second  
CAM block.

32. A method of operating a content addressable memory  
(CAM) device having a plurality of CAM blocks each  
including an array of CAM cells, comprising:

storing data of a first class type in a first CAM block;

storing data of a second class type in a second CAM  
block;

receiving comparand data associated with the first class  
type; and

selectively comparing the comparand data only with data  
of the first class type of the first CAM block.

33. The method of claim 32, wherein the selectively  
comparing comprises:

enabling the comparand data to be provided to the first  
CAM block; and

disabling the comparand data from being provided to the  
second CAM block.

\* \* \* \* \*